



合肥学院
HEFEI UNIVERSITY



Optimization Algorithms

5. Stochastic Hill Climbing

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Introduction



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- Each search step consists of creating an entirely new, entirely random candidate solution.
- Each search step is thus independent of all prior steps.
- **Is this a good idea?**
- Probably not.
- In almost all practical scenarios, good solutions are somewhat similar to other good solutions.
- In other words, every good solution we see is some useful information.

Basic Idea

- So how we can make use of the information we have seen during the search?

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- So how we can make use of the information we have seen during the search?
- Instead of generating a completely random new candidate solution in each step. . .
- . . . why can't we try to iteratively improve the best solution we have seen so far?

Algorithm Concept



Stochastic Hill Climbing

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 3. if it is better, it becomes the new best-so-far solution (if it is not better, discard it).

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 1. create random initial solution
 2. make a modified copy of best-so-far solution
 3. if it is better, it becomes the new best-so-far solution (if it is not better, discard it).
 4. go back to 2. (until the time is up)

Implementation of the Stochastic Hill Climber

```
package aitoa.algorithms;

public class HillClimber<X, Y> {
// unnecessary stuff omitted here...
//
//
//
//

//
//

//
//
//
//
//
//
//
//
//
//
}
```


Implementation of the Stochastic Hill Climber

```
package aitoa.algorithms;

public class HillClimber<X, Y> extends Metaheuristic1<X, Y> {
    // unnecessary stuff omitted here...
    public void solve(IBlackBoxProcess<X, Y> process) {
        X xCur = process.getSearchSpace().create();
        X xBest = process.getSearchSpace().create();
    }
}
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    public void solve(IBlackBoxProcess<X, Y> process) {
        X xCur = process.getSearchSpace().create();
        X xBest = process.getSearchSpace().create();
        Random random = process.getRandom();

        this.nullary.apply(xBest, random);
        double fBest = process.evaluate(xBest);

        //
        this.unary.apply(xBest, xCur, random);
        //
        //
        //
        //
        //
        //
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        //
        this.unary.apply(xBest, xCur, random);
        double fCur = process.evaluate(xCur);

        //
        //
        //
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}
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Implementation of the Stochastic Hill Climber

```
package aitoa.algorithms;

public class HillClimber<X, Y> extends Metaheuristic1<X, Y> {
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    public void solve(IColorProcess<X, Y> process) {
        X xCur = process.getSearchSpace().create();
        X xBest = process.getSearchSpace().create();
        Random random = process.getRandom();

        this.nullary.apply(xBest, random);
        double fBest = process.evaluate(xBest);

        //
        this.unary.apply(xBest, xCur, random);
        double fCur = process.evaluate(xCur);
        if (fCur < fBest) {
            //
            //
        }
        //
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        double fCur = process.evaluate(xCur);
        if (fCur < fBest) {
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        }
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        this.nullary.apply(xBest, random);
        double fBest = process.evaluate(xBest);

    //
        this.unary.apply(xBest, xCur, random);
        double fCur = process.evaluate(xCur);
        if (fCur < fBest) {
            fBest = fCur;
            process.getSearchSpace().copy(xCur, xBest);
        }
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        Random random = process.getRandom();

        this.nullary.apply(xBest, random);
        double fBest = process.evaluate(xBest);

        while (!process.shouldTerminate()) {
            this.unary.apply(xBest, xCur, random);
            double fCur = process.evaluate(xCur);
            if (fCur < fBest) {
                fBest = fCur;
                process.getSearchSpace().copy(xCur, xBest);
            }
        }
    }
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Causality

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- Causality means that small changes in the features of an object (or candidate solution) also lead to small changes in its behavior (or objective value).
- If an optimization problem exhibits causality, then there should be good solutions that are similar to other good solutions.
- The idea is that if we have a good candidate solution, then there may exist similar solutions which are better.
- We hope to find one of them and then continue trying to do the same from there.

Ingredient: Unary Search Operator



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- It should ideally be randomized, i.e., applying it twice to the same point x should yield different results.

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```
package aitoa.structure;  
  
public interface IUnarySearchOperator<X> {  
    void apply(X x, X dest, Random random);  
}
```

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- How can we implement this for our JSSP scenario?
- **Easy: Just swap two (different) job IDs in the string!**
- Since the numbers of occurrences of the IDs will not change, the new strings will be valid.

Example for our 1swap Operator



(2,0,1,0,1,1,2,3,2,3,
2,0,0,1,3,3,2,3,1,0)

Example for our 1swap Operator

X

(2,0,1,0,1,1,2,3,2,3,
2,0,0,1,3,3,2,3,1,0)



Y

Example for our 1swap Operator

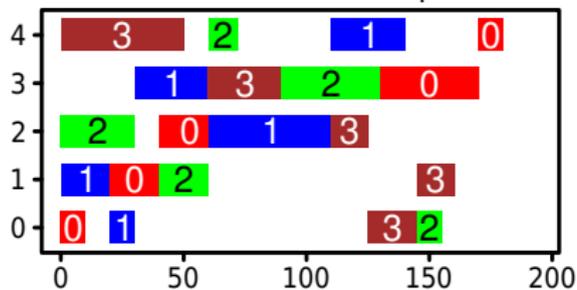
X

(2,0,1,0,1,1,2,3,2,3,
2,0,0,1,3,3,2,3,1,0)



makespan: 180

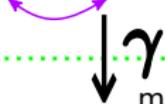
Y



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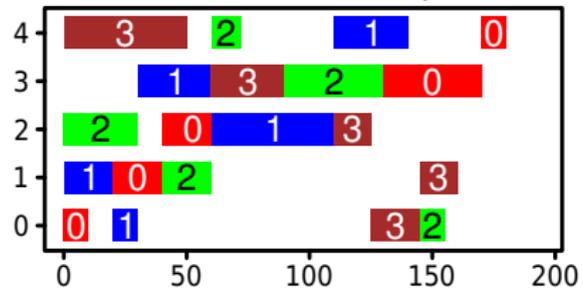
X

(2,0,1,0,1,1,2,3,2,3,
2,0,0,1,3,3,2,3,1,0)



makespan: 180

Y

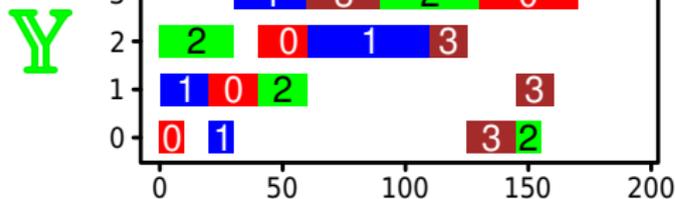


Example for our 1swap Operator

X $(2,0,1,0,1,1,2,3,2,3, 2,0,0,1,3,3,2,3,1,0) \xrightarrow{1\text{swap}}$ $(2,0,1,0,1,1,2,3,2,3, 2,0,3,1,3,0,2,3,1,0)$



makespan: 180



Example for our 1swap Operator

X

(2,0,1,0,1,1,2,3,2,3,
2,0,0,1,3,3,2,3,1,0)

1swap

(2,0,1,0,1,1,2,3,2,3,
2,0,3,1,3,0,2,3,1,0)

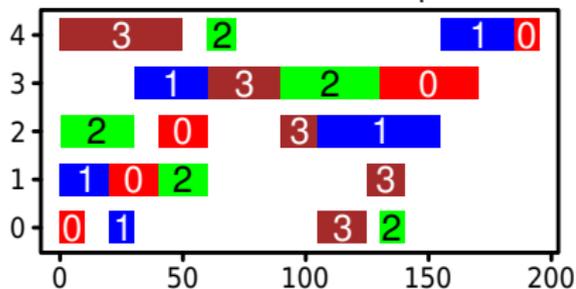
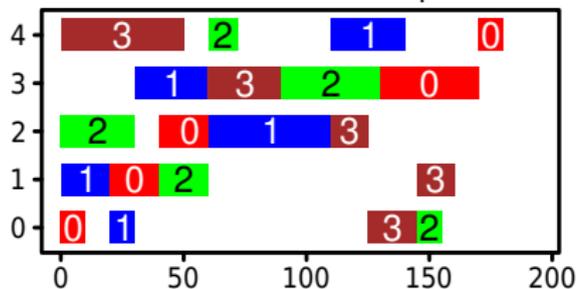
γ

makespan: 180

γ

makespan: 195

Y



Example for our 1swap Operator

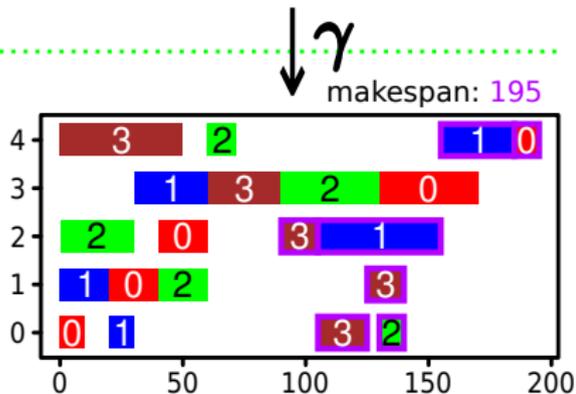
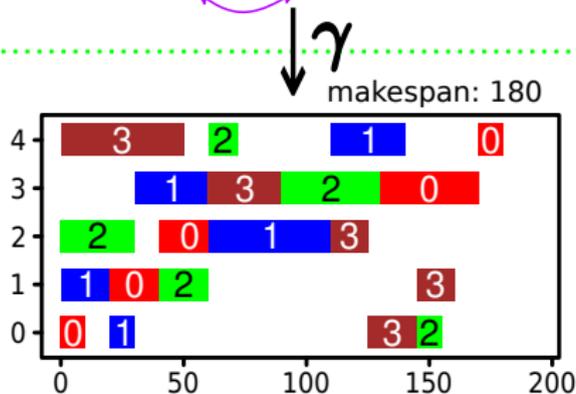
X

(2,0,1,0,1,1,2,3,2,3,
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1swap →

(2,0,1,0,1,1,2,3,2,3,
2,0,3,1,3,0,2,3,1,0)

Y



Example for our 1swap Operator

X

(2,0,1,0,1,1,2,3,2,3,
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← 1swap

(2,0,1,0,1,1,2,3,2,3,
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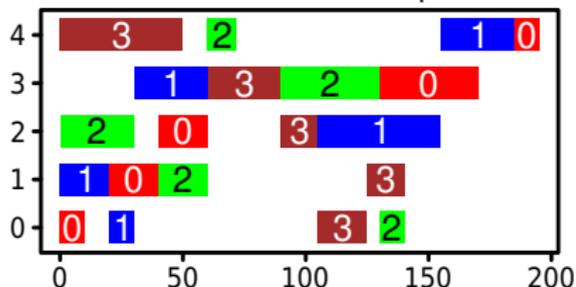
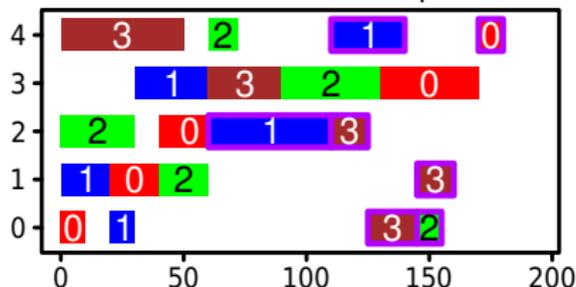


makespan: 180



makespan: 195

Y




```
package aitoa.examples.jssp;

public class JSSPUnaryOperator1Swap implements
    IUnarySearchOperator<int[]> {
    // unnecessary stuff omitted here...
    public void apply(int[] x, int[] dest, Random random) {
    // copy the source point in search space to the dest
        System.arraycopy(x, 0, dest, 0, x.length);

    // choose the index of the first operation to swap
        int i = random.nextInt(dest.length);
        int jobI = dest[i]; // remember job id

    //
    //
    //
    //
    //
    //
    //
    //
    //
    //
    }
}
```

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        System.arraycopy(x, 0, dest, 0, x.length);

    // choose the index of the first operation to swap
        int i = random.nextInt(dest.length);
        int jobI = dest[i]; // remember job id

    // choose index of second operation to swap
        int j = random.nextInt(dest.length);
        int jobJ = dest[j];

    //
    //
    //
    //
    //
    //
    }
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        int j = random.nextInt(dest.length);
        int jobJ = dest[j];

    //
        dest[i] = jobJ;
        dest[j] = jobI; // then we swap the values

    //
    //
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    // choose the index of the first operation to swap
        int i = random.nextInt(dest.length);
        int jobI = dest[i]; // remember job id

    // choose index of second operation to swap
        int j = random.nextInt(dest.length);
        int jobJ = dest[j];
        if (jobI != jobJ) { // we found two locations with two
            dest[i] = jobJ; // different values
            dest[j] = jobI; // then we swap the values
        }
    }
}
```

```
package aitoa.examples.jssp;

public class JSSPUnaryOperator1Swap implements
    IUnarySearchOperator<int[]> {
    // unnecessary stuff omitted here...
    public void apply(int[] x, int[] dest, Random random) {
    // copy the source point in search space to the dest
        System.arraycopy(x, 0, dest, 0, x.length);

    // choose the index of the first sub-job to swap
        int i = random.nextInt(dest.length);
        int jobI = dest[i]; // remember job id

        for (;;) { // try to find a location j with a different job
            int j = random.nextInt(dest.length);
            int jobJ = dest[j];
            if (jobI != jobJ) { // we found two locations with two
                dest[i] = jobJ; // different values
                dest[j] = jobI; // then we swap the values
                return; // and are done
            }
        }
    }
}
```

Experiment and Analysis



So what do we get?

- I execute the program 101 times for each of the instances abz7, la24, swv15, and yn4

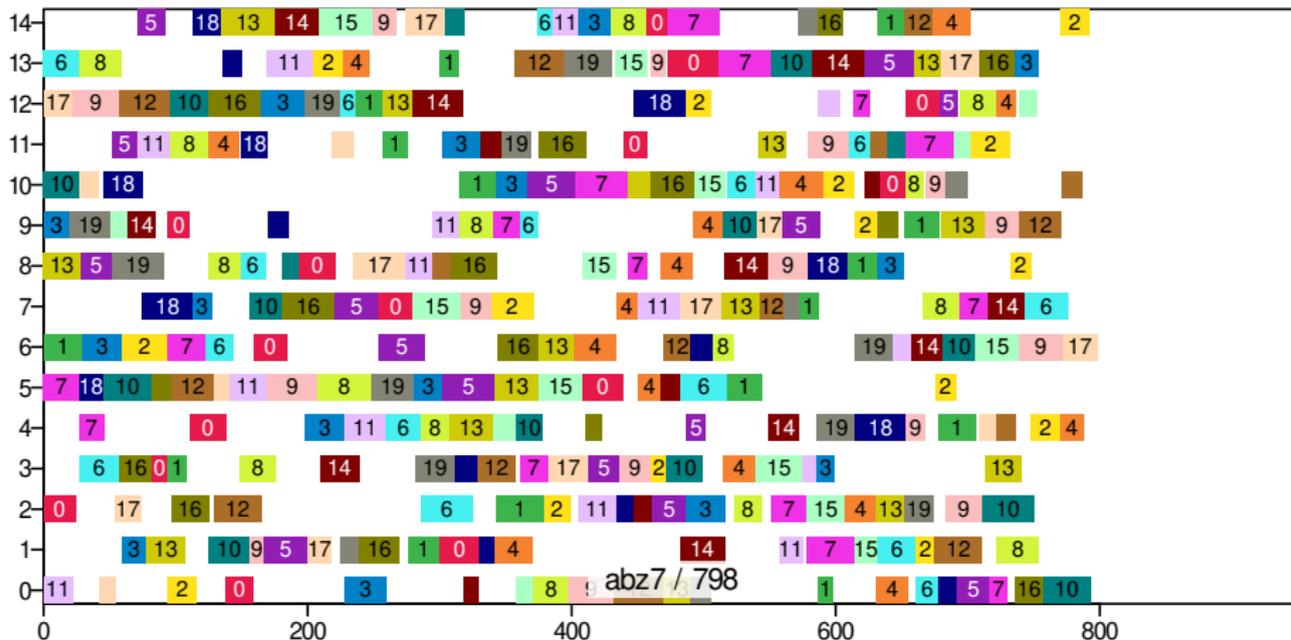
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\mathcal{I}	algo	makespan				last improvement	
		best	mean	med	sd	med(t)	med(FEs)
abz7	rs	895	947	949	12	85s	6'512'505
	hc_1swap	717	800	798	28	0s	16'978
la24	rs	1153	1206	1208	15	82s	15'902'911
	hc_1swap	999	1095	1086	56	0s	6'612
swv15	rs	4988	5166	5172	50	87s	5'559'124
	hc_1swap	3837	4108	4108	137	1s	104'598
yn4	rs	1460	1498	1499	15	76s	4'814'914
	hc_1swap	1109	1222	1220	48	0s	31'789

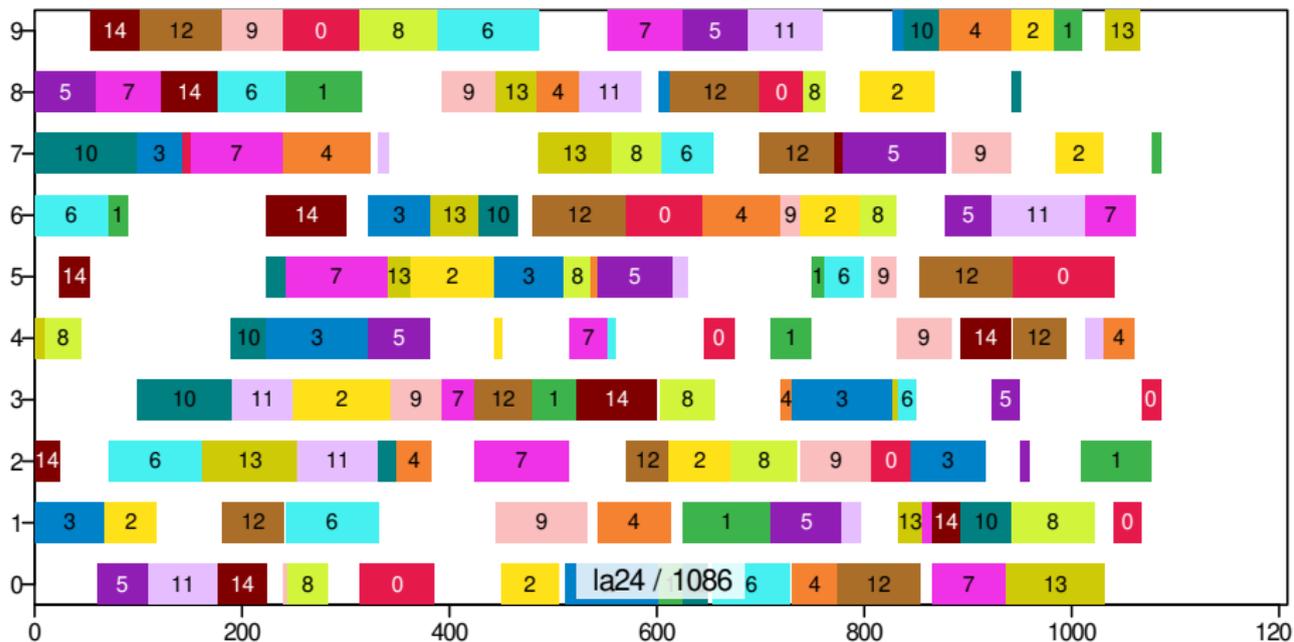
So what do we get?

hc_1swap: median result of 3 min of hill climber



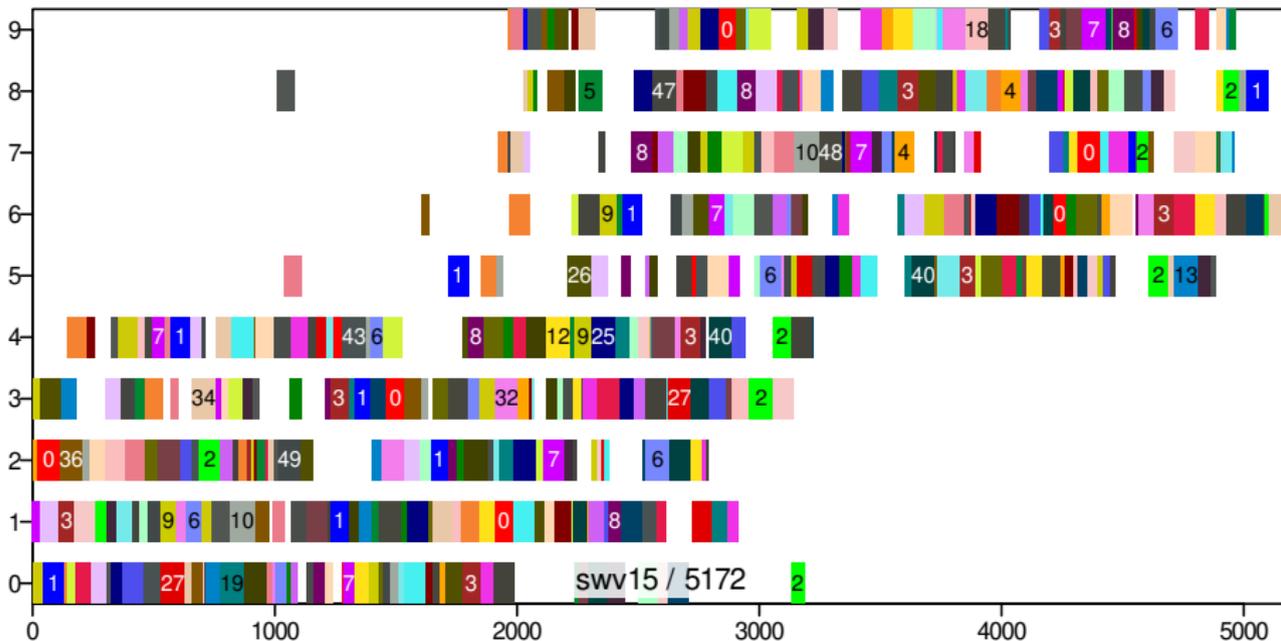
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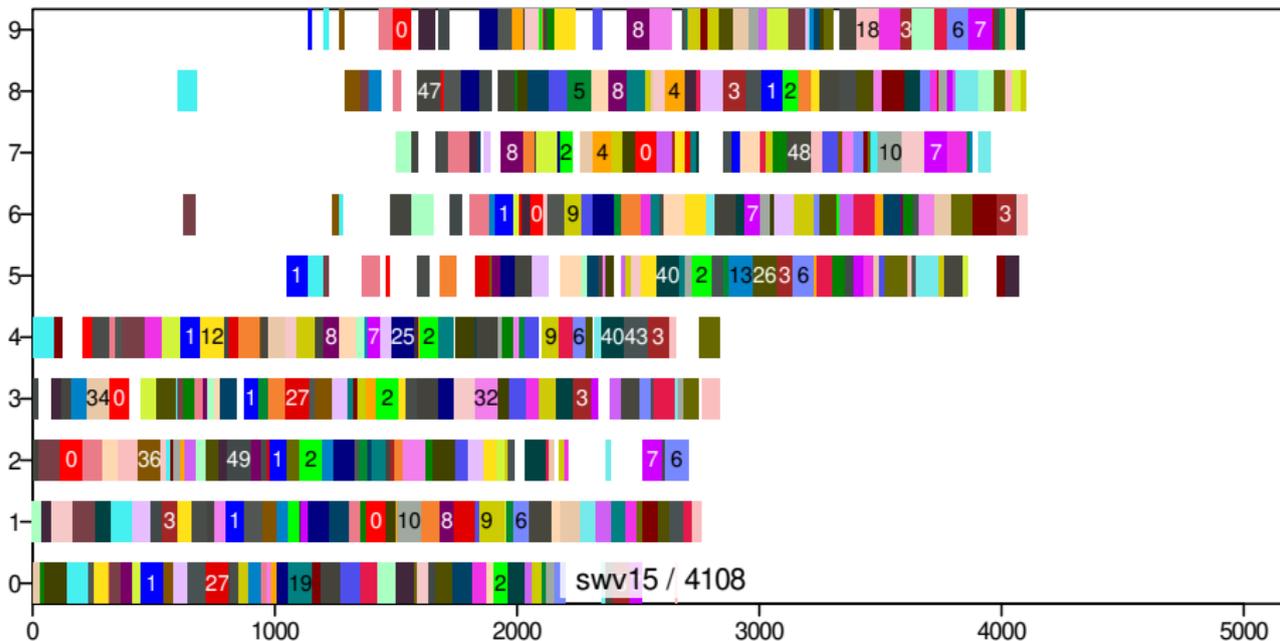
So what do we get?

rs: median result of 3 min of random sampling



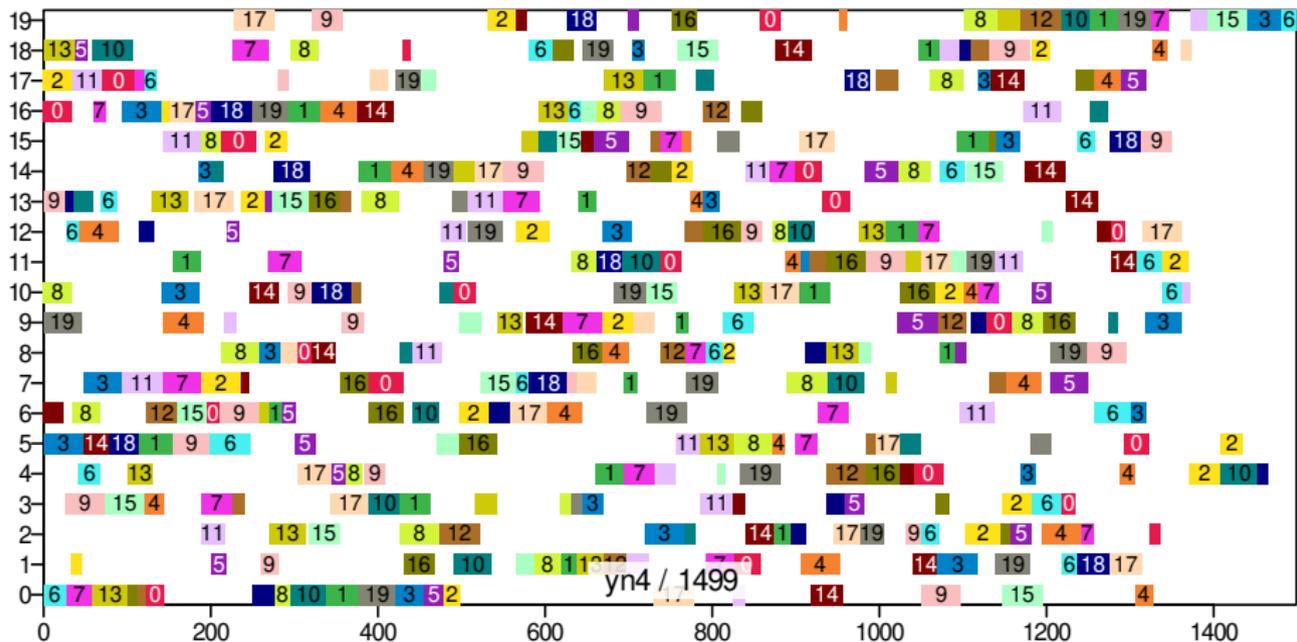
So what do we get?

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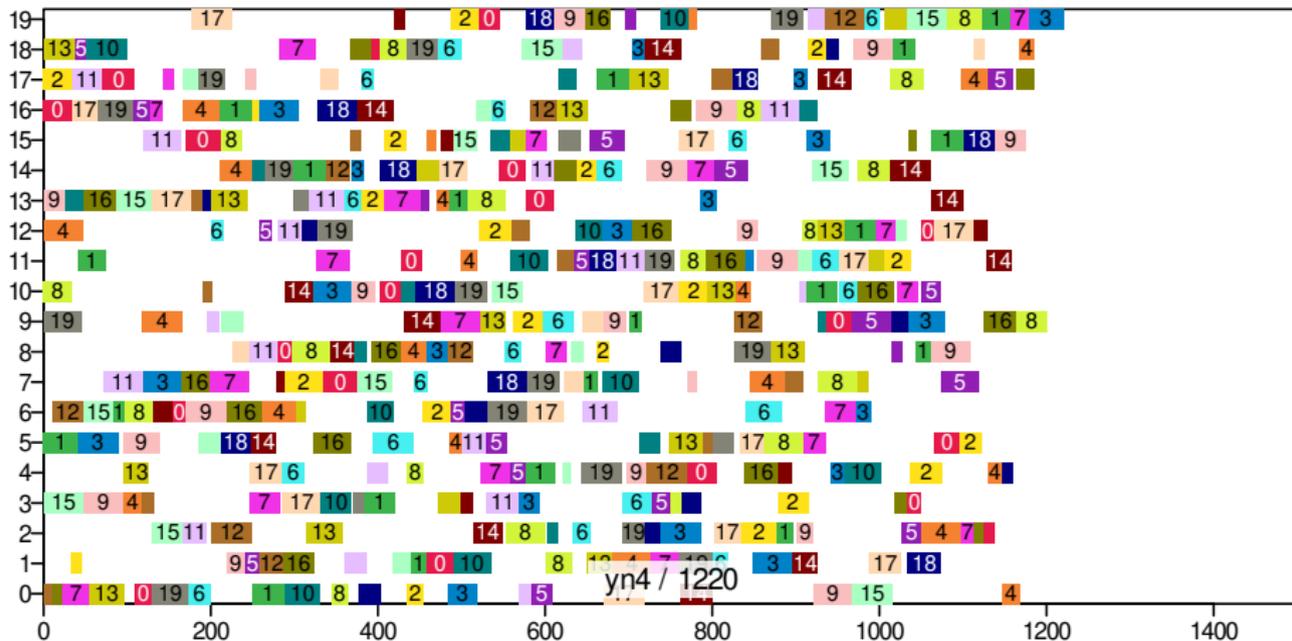
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hc_1swap: median result of 3 min of hill climber

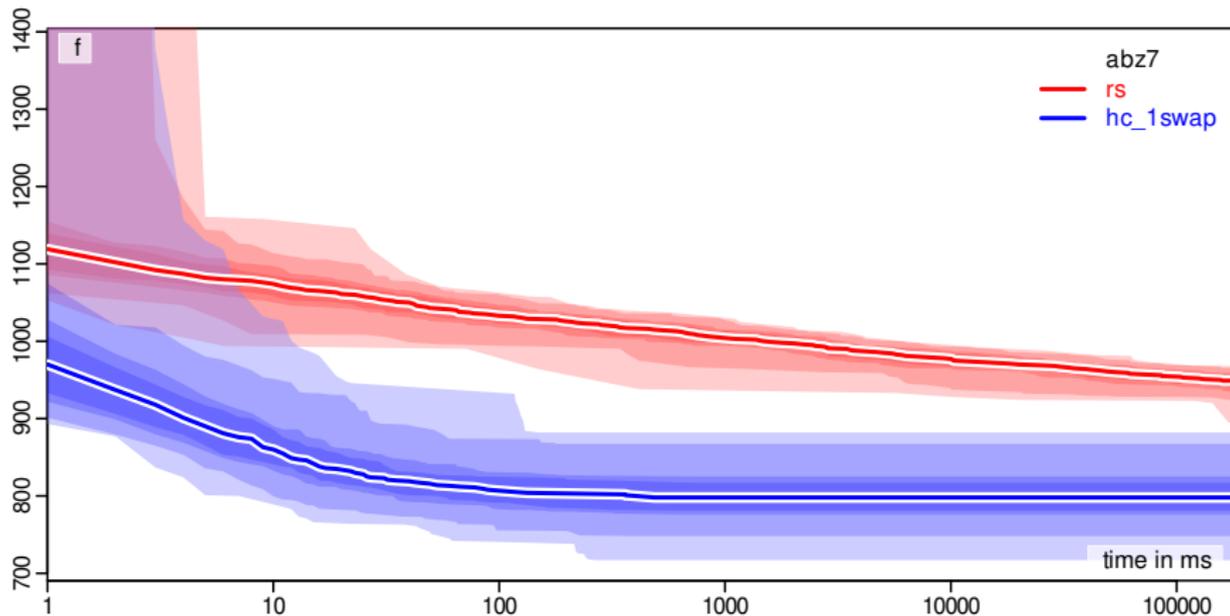


Progress over Time

What progress does the algorithm make over time?

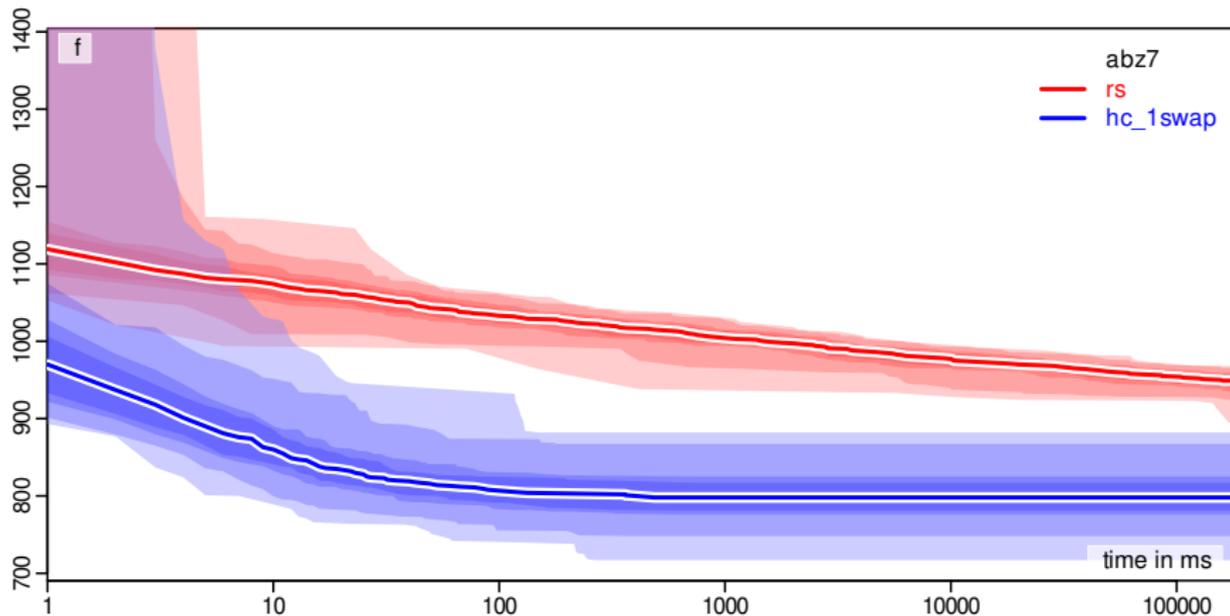
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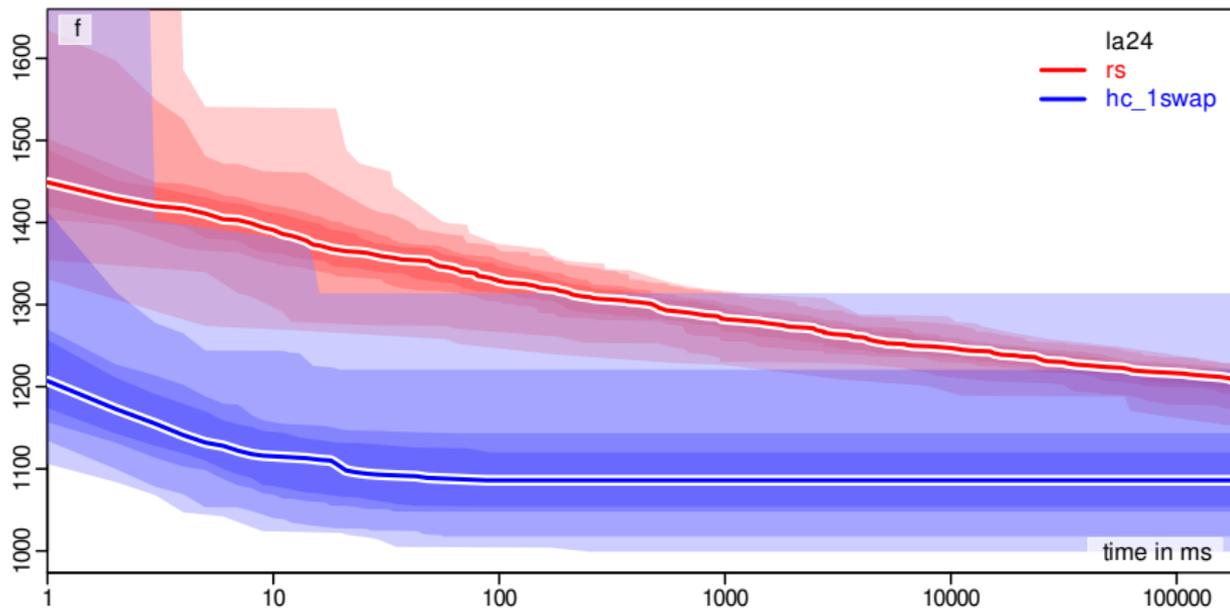
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First we have much progress...

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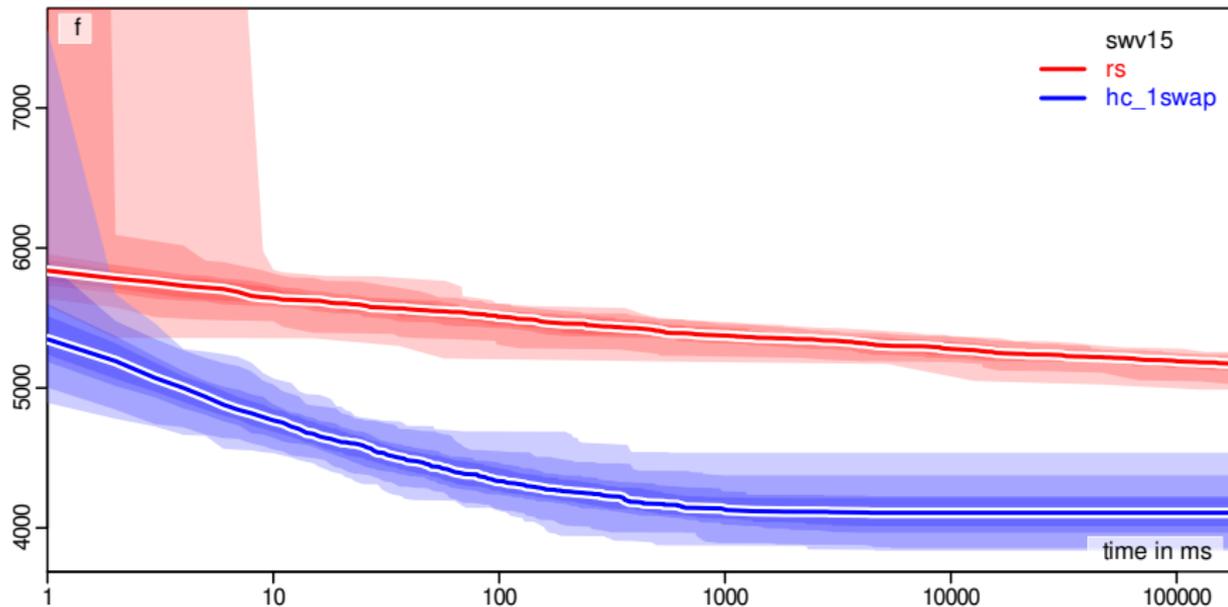
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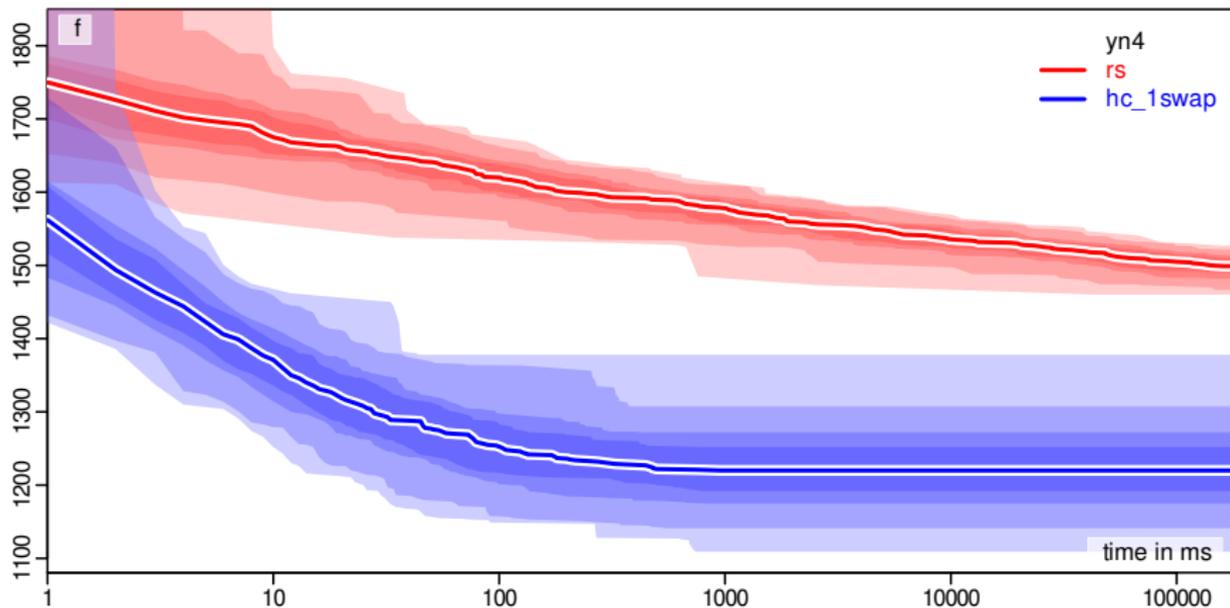
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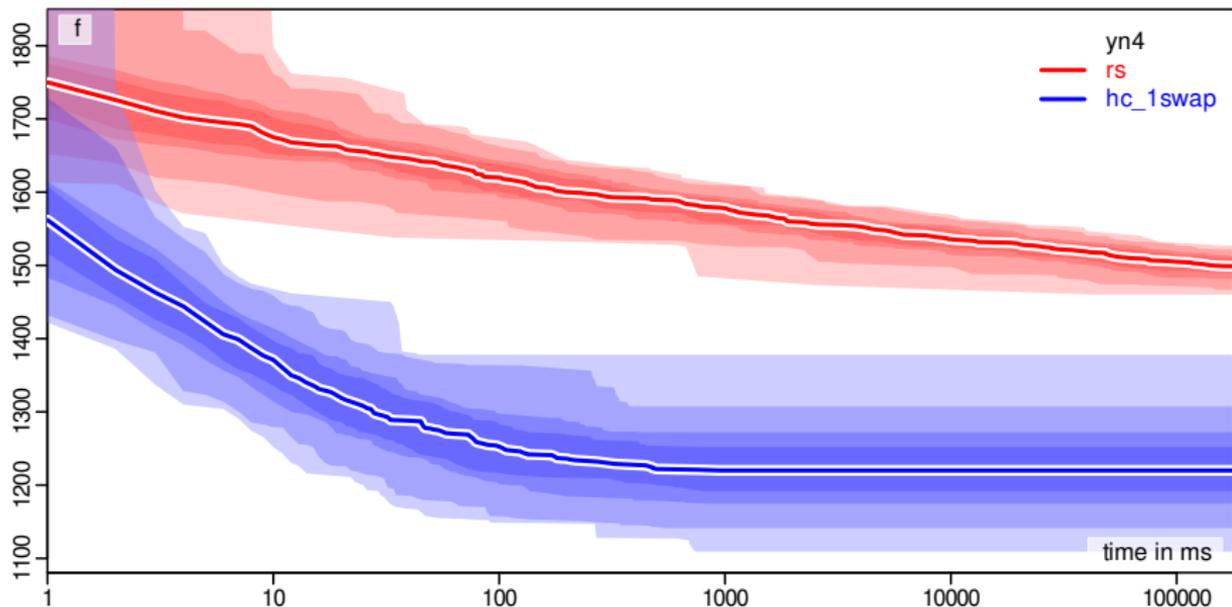
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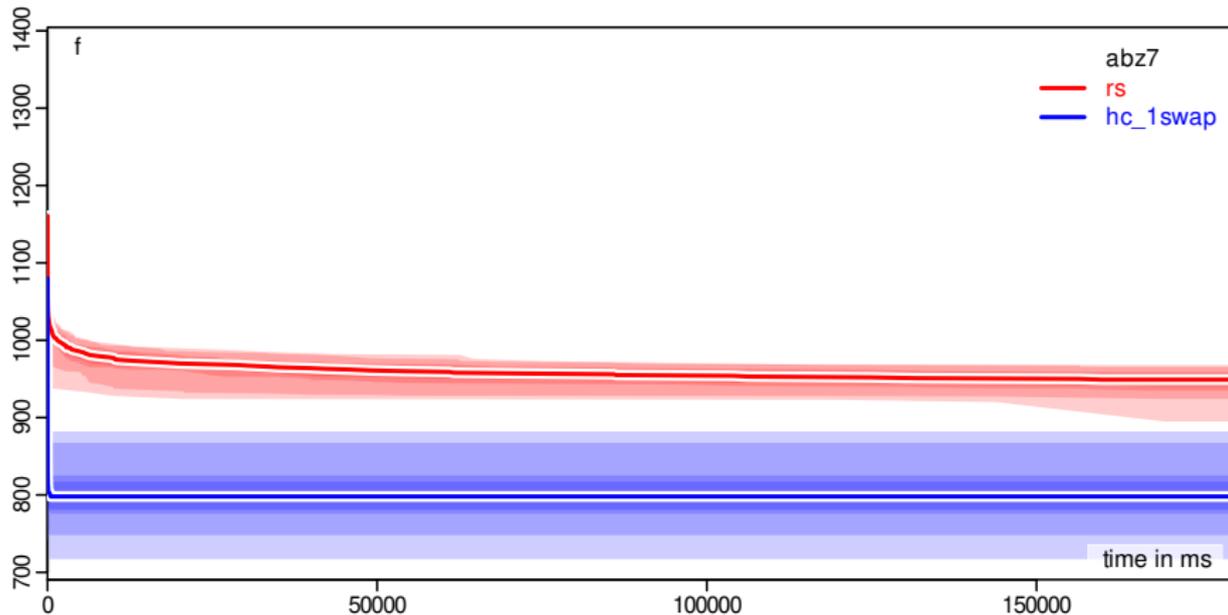
First we have much progress...
... but then the hill climber stagnates!

But we waste time...

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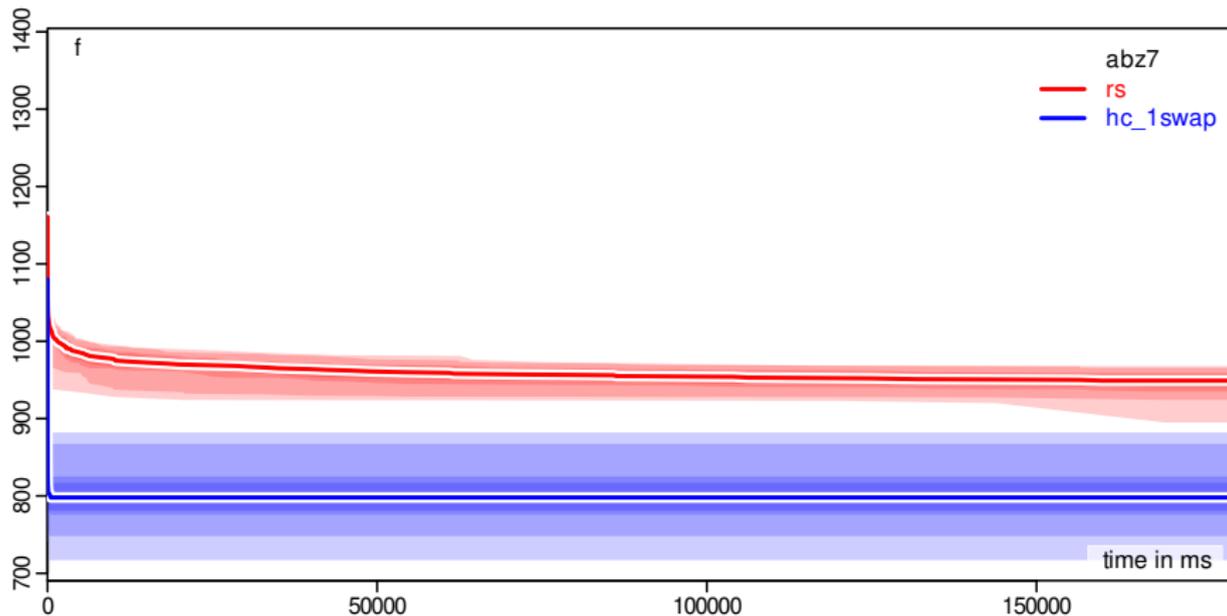
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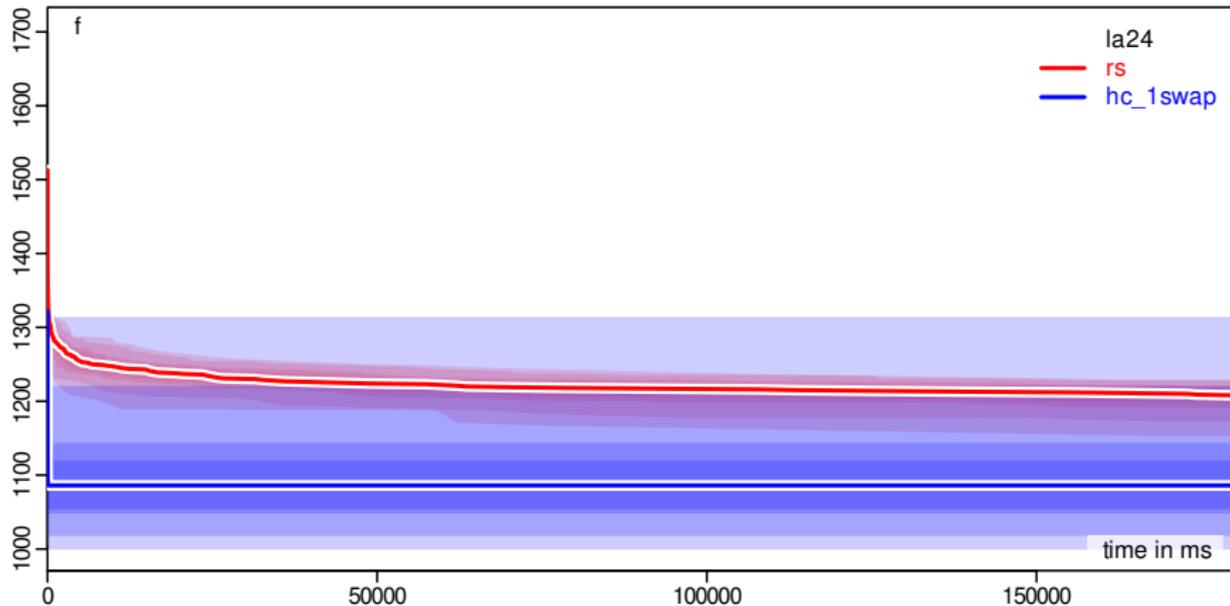
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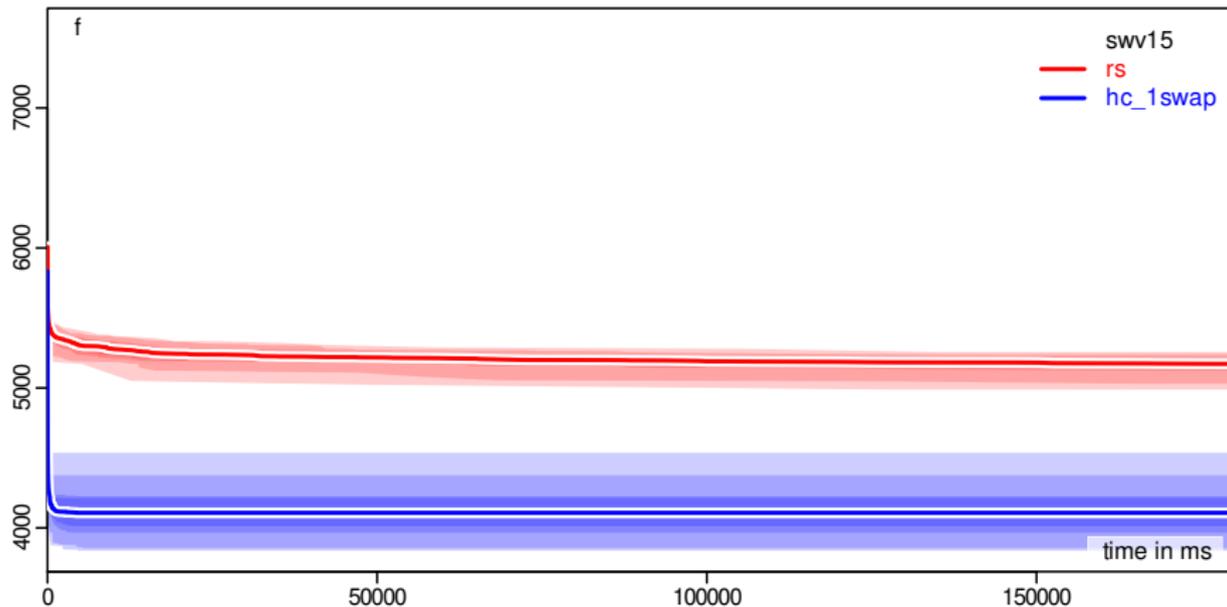
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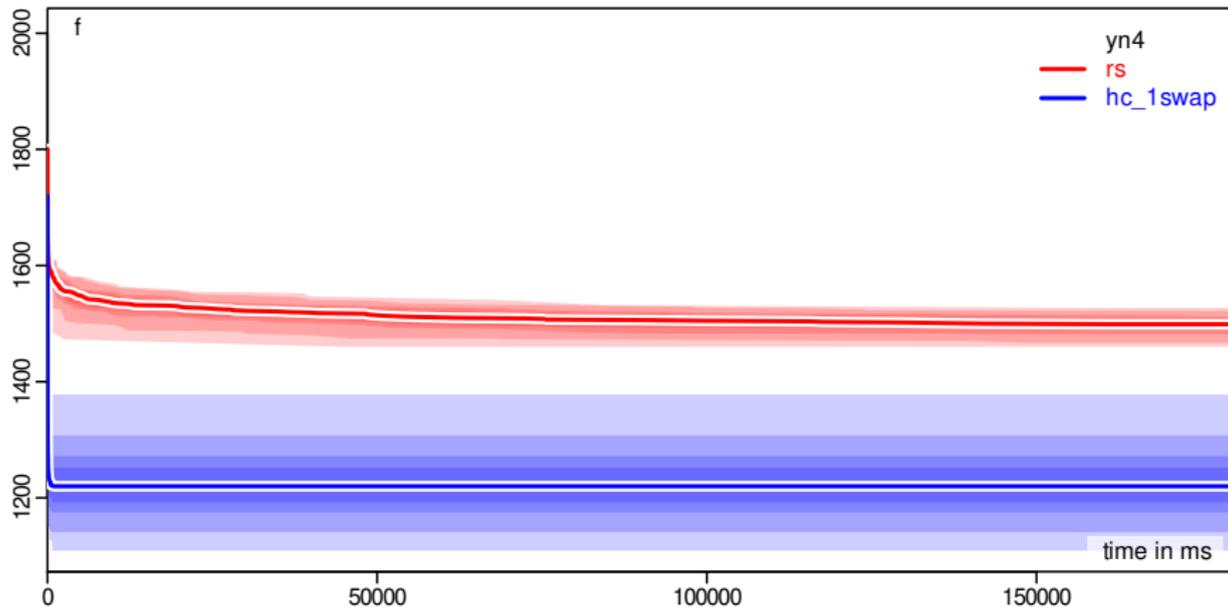
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- We have three minutes but after about 1 second, our hc_1swap algorithm stops improving!

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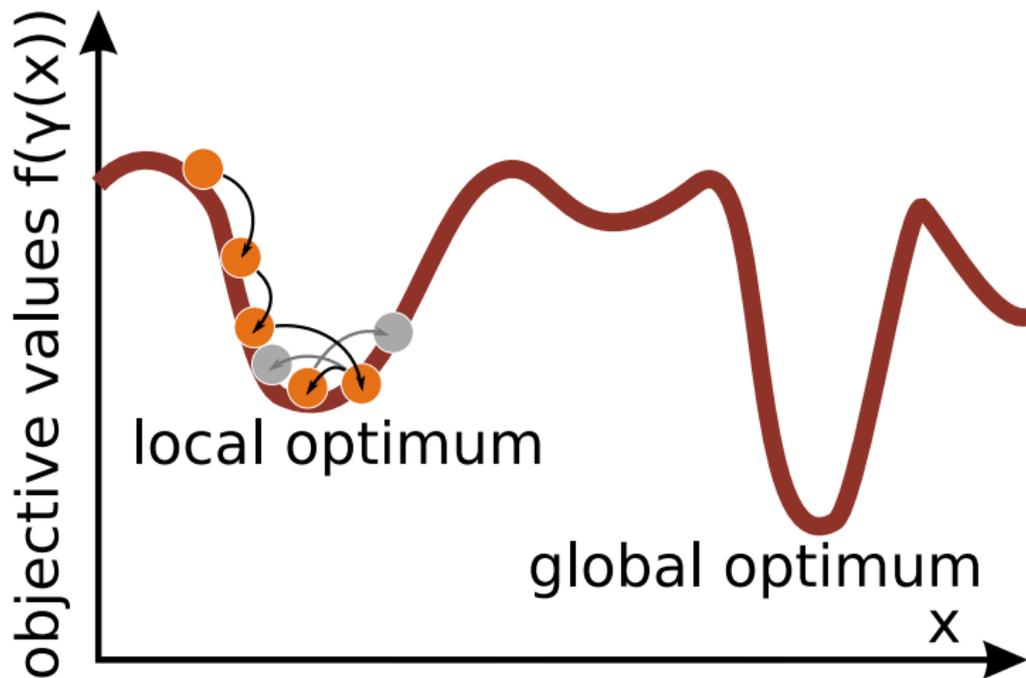
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- This is called **Premature Convergence**.^{8 9}

Premature Convergence



Improved Algorithm Concept 1



Stochastic Hill Climber with Restarts

- Idea: We have seen that the results of the hill climber exhibit a relatively **high standard deviation**.

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- Idea: If we did not make any progress for a number L of algorithm steps, we simply restart at a new random solution.
- Of course, we will always remember the overall best solution we ever had (in another variable).

Stochastic Hill Climbing Algorithm with Restarts

```
package aitoa.algorithms;

public class HillClimberWithRestarts<X, Y> extends Metaheuristic1<X, Y> {
    // unnecessary stuff omitted here...
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Experiment and Analysis



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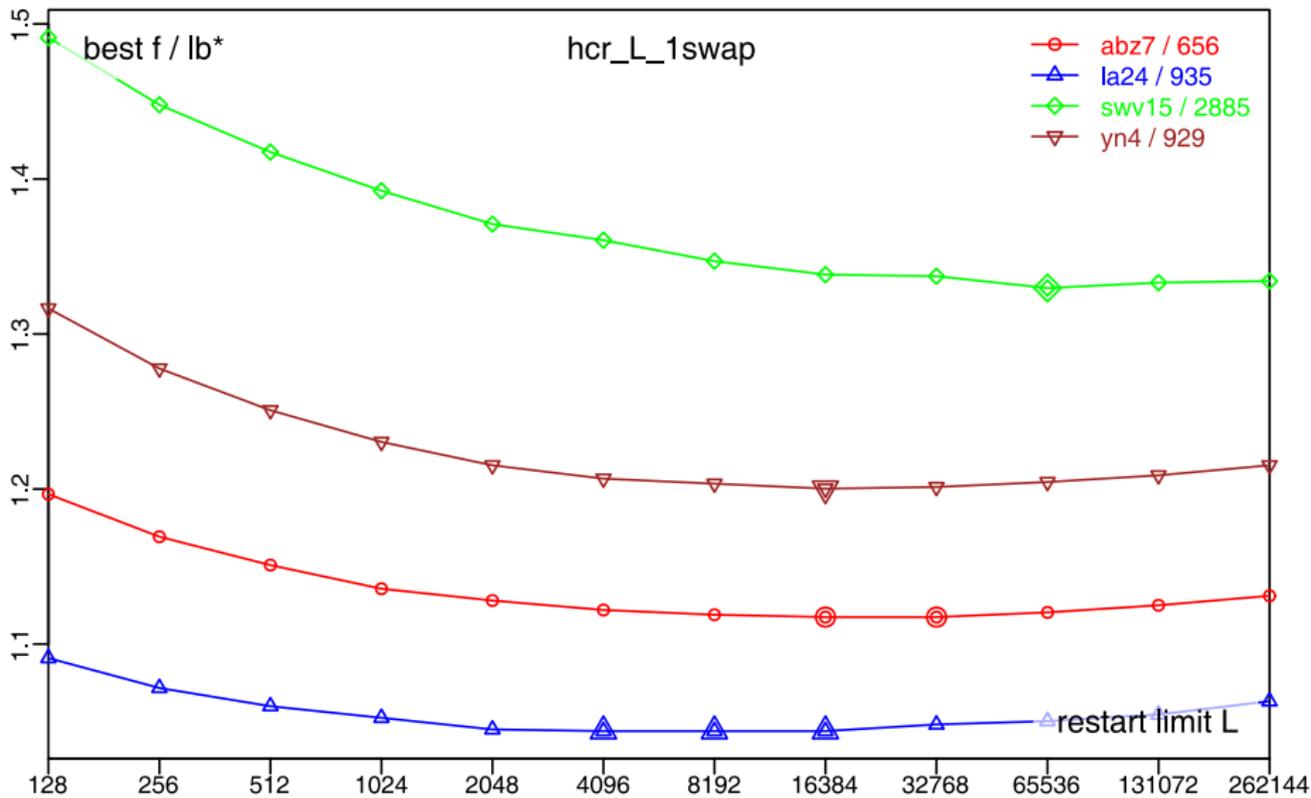
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- $L = 2^{14} = 16'384$ seems to be a reasonable choice.

So what do we get?

- I execute the program 101 times for each of the instances abz7, la24, swv15, and yn4

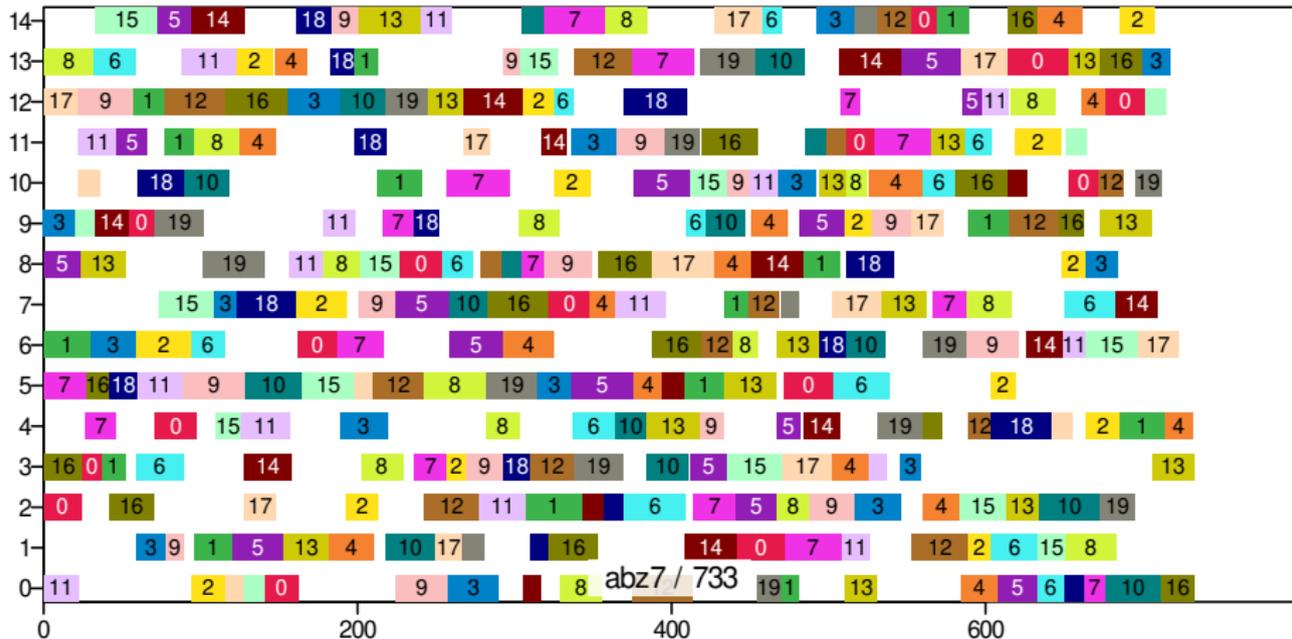
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	hc_1swap	999	1095	1086	56	0s	6'612
	hcr_16384_1swap	953	976	976	7	80s	34'437'999
swv15	rs	4988	5166	5172	50	87s	5'559'124
	hc_1swap	3837	4108	4108	137	1s	104'598
	hcr_16384_1swap	3752	3859	3861	42	92s	11'756'497
yn4	rs	1460	1498	1499	15	76s	4'814'914
	hc_1swap	1109	1222	1220	48	0s	31'789
	hcr_16384_1swap	1081	1115	1115	11	91s	14'804'358

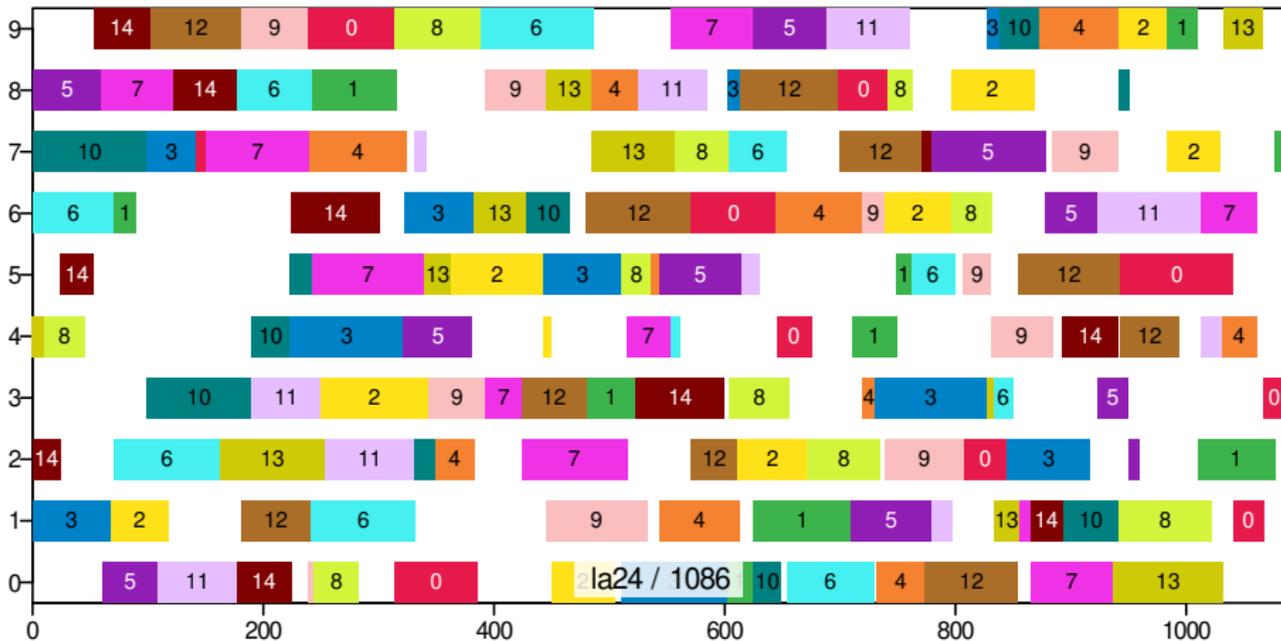
So what do we get?

hcr_16384_1swap: median result of 3 min of hill climber which restarts after $L = 16'384$ search steps without improvement



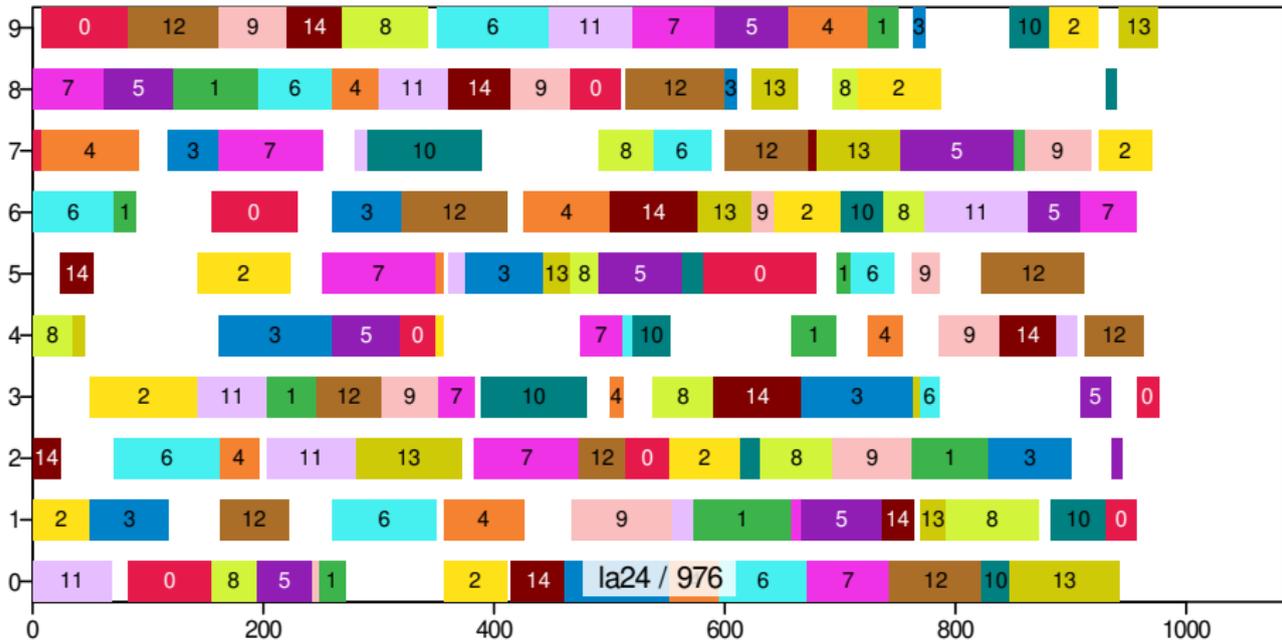
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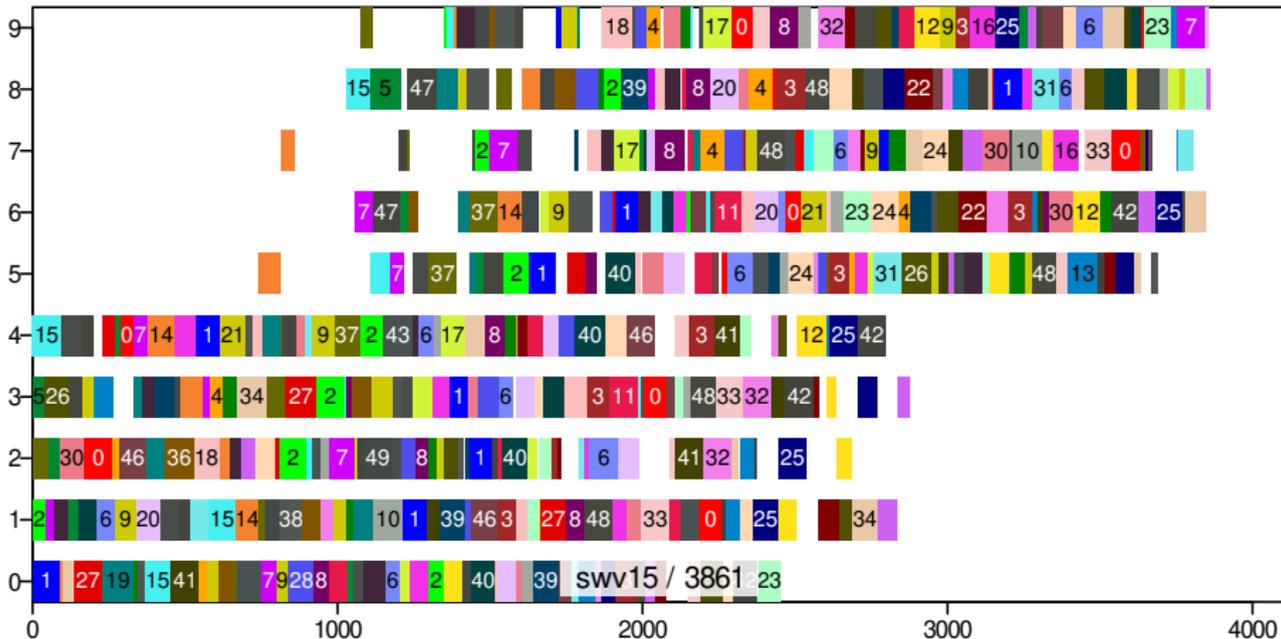
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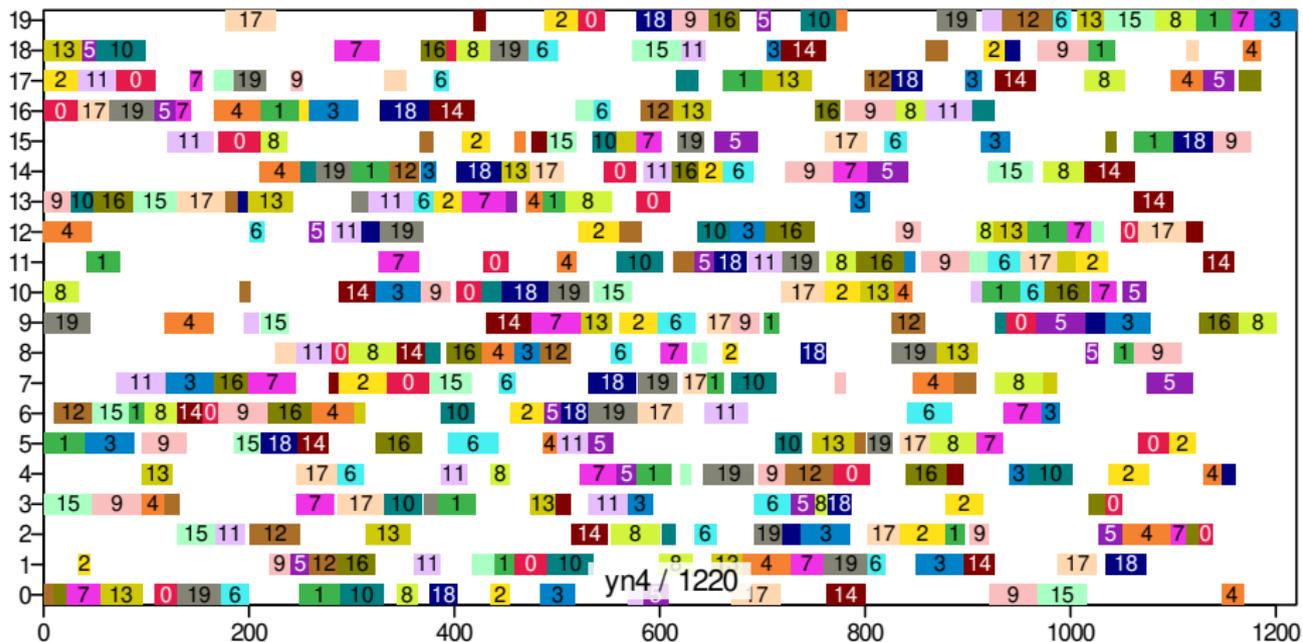
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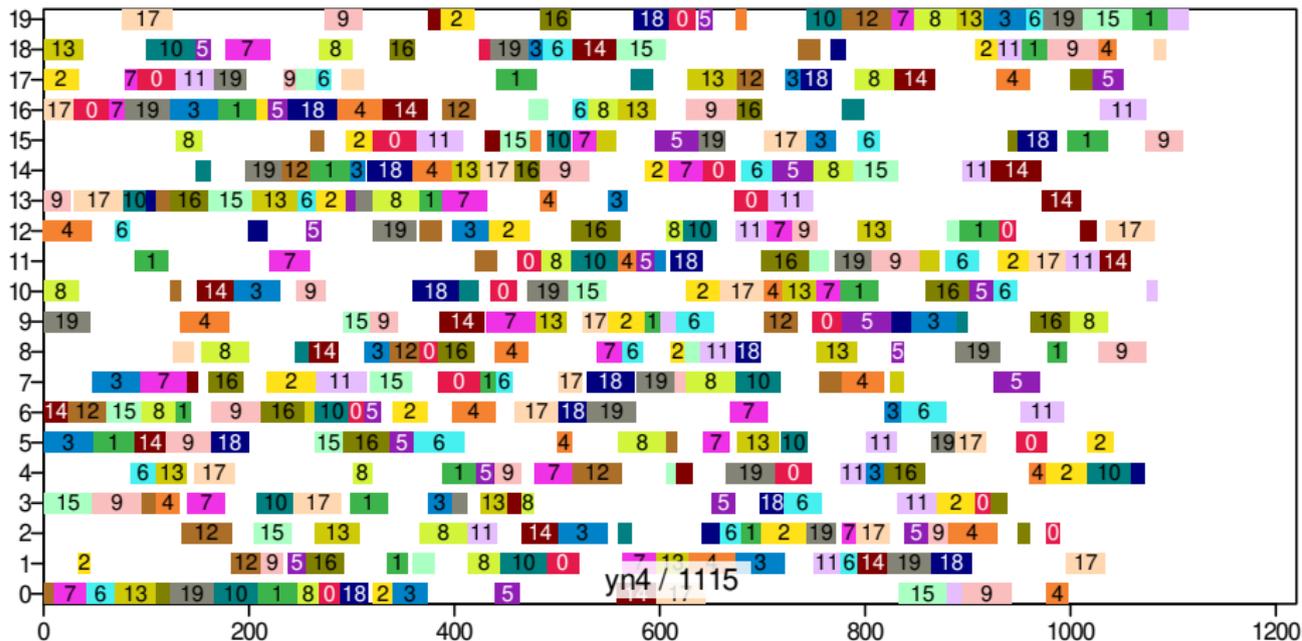
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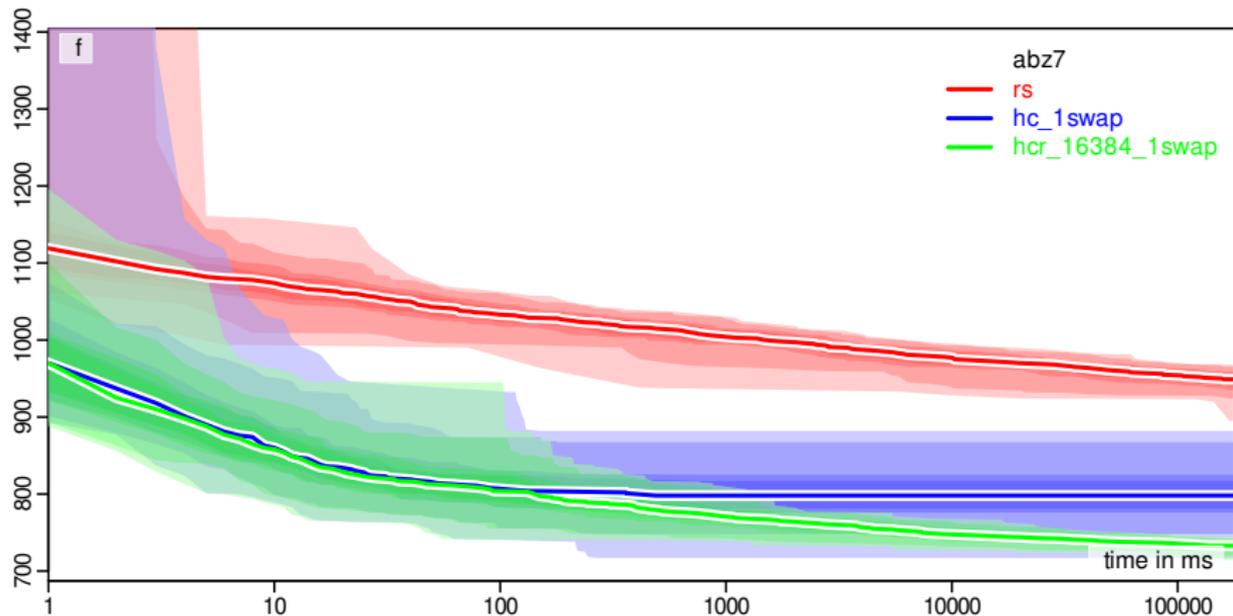
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Progress over Time

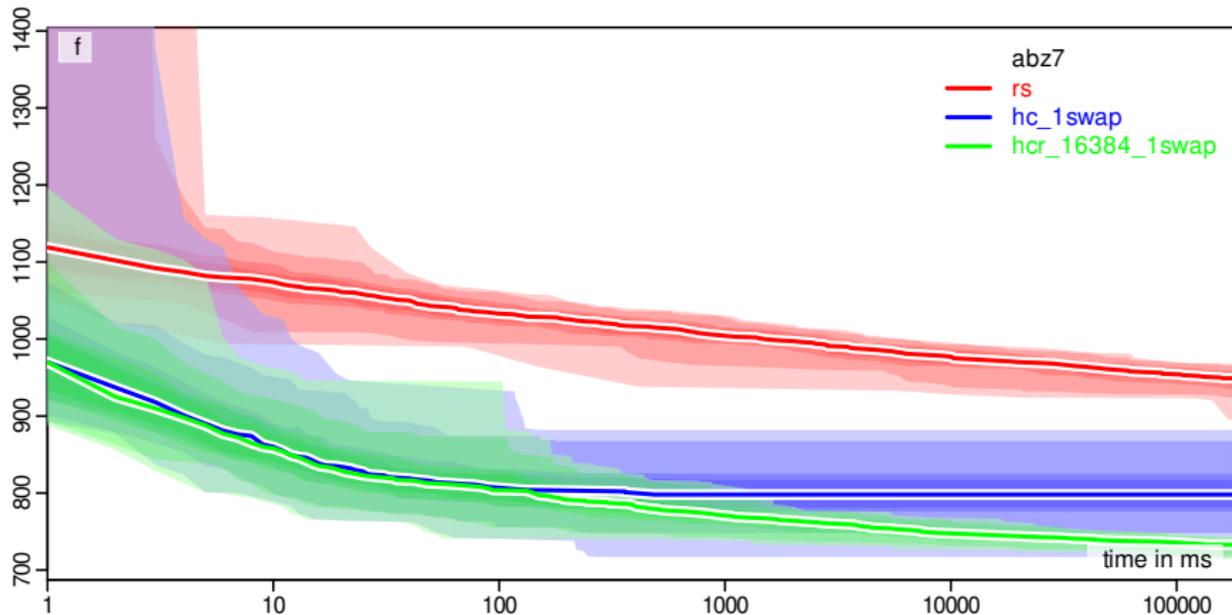
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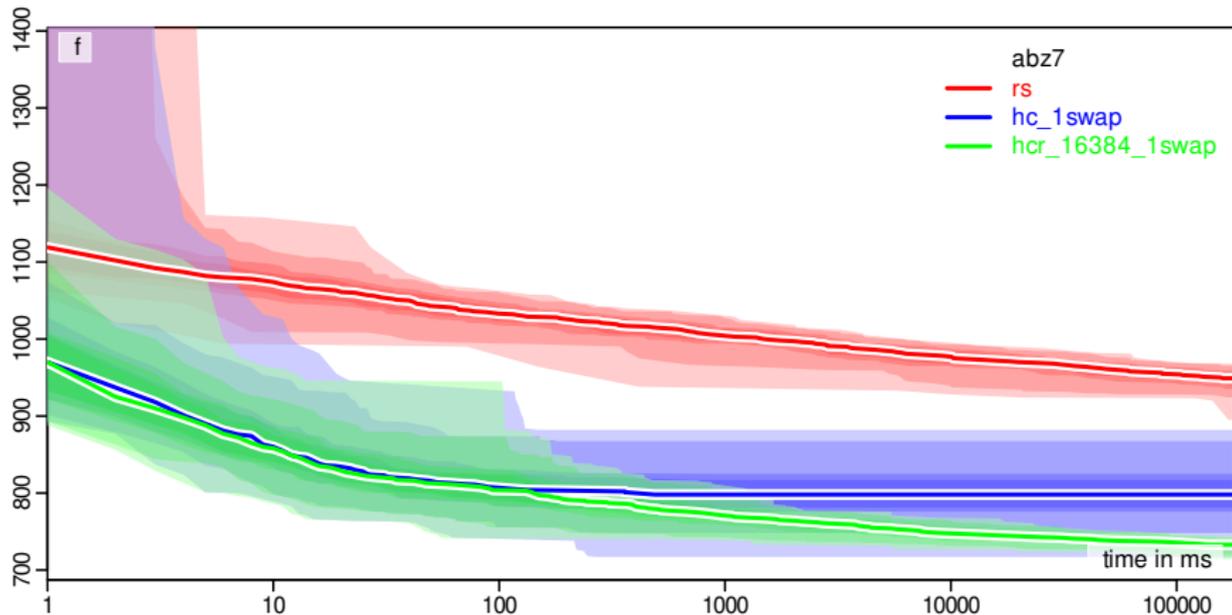
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What progress does the algorithm make over time?

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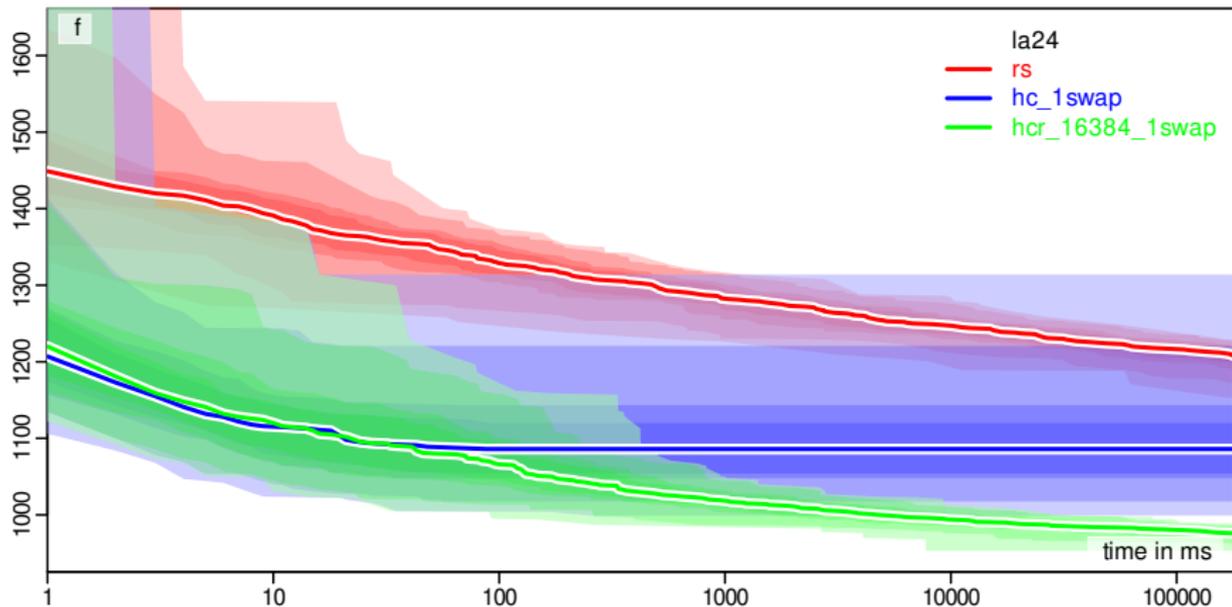
Progress over Time



What progress does the algorithm make over time?

- First it behaves like the normal hill climber
- But it keeps improving.

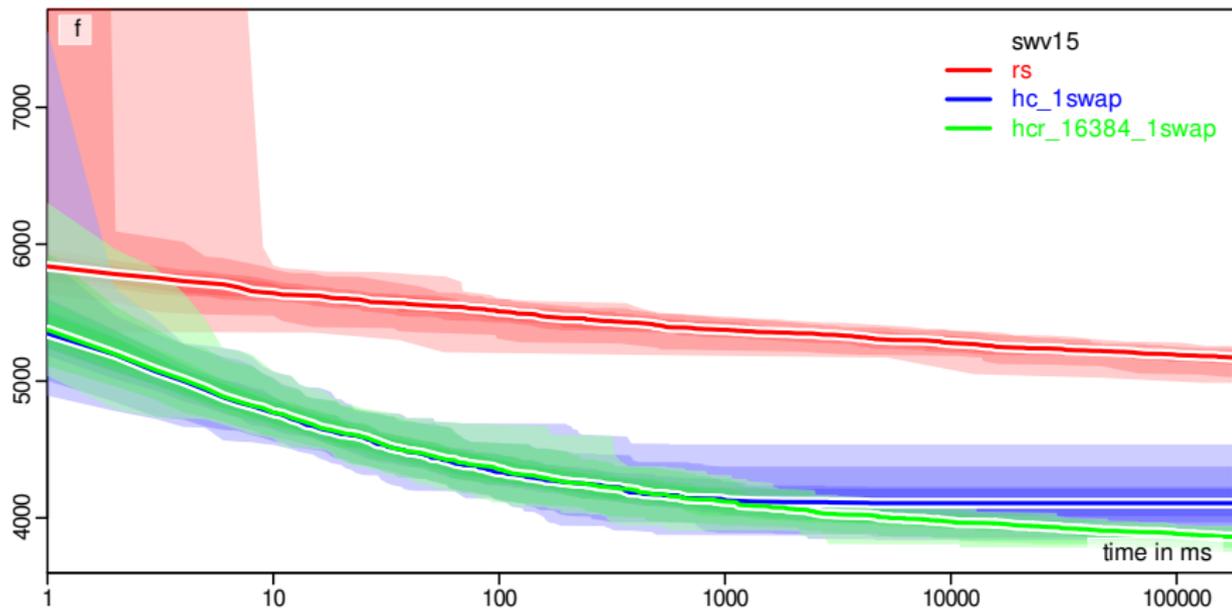
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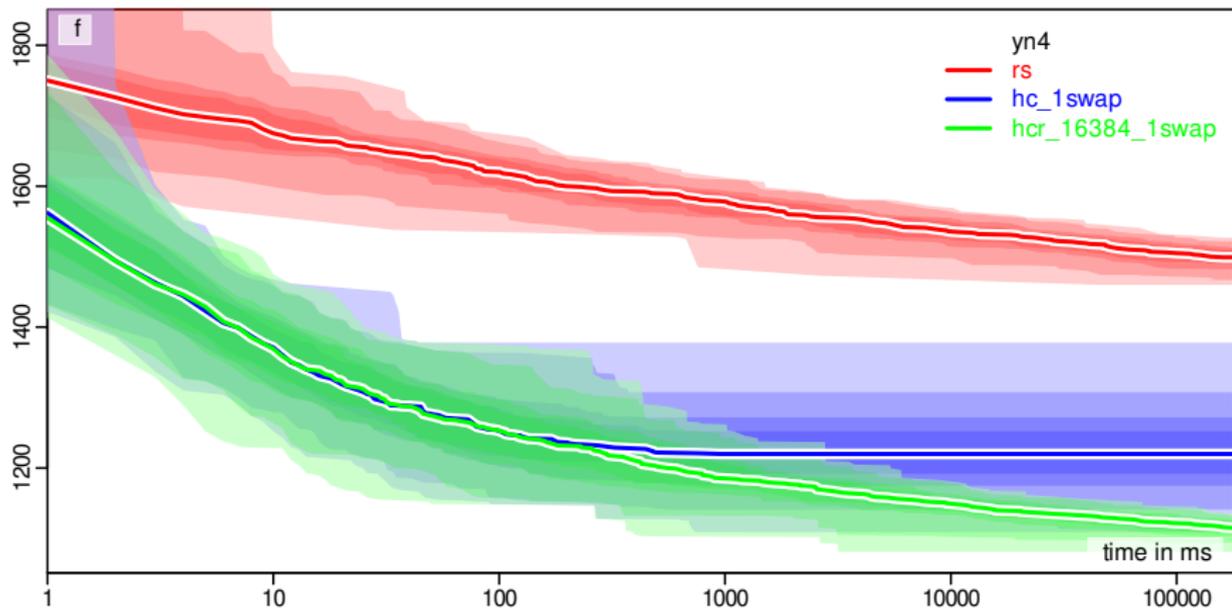
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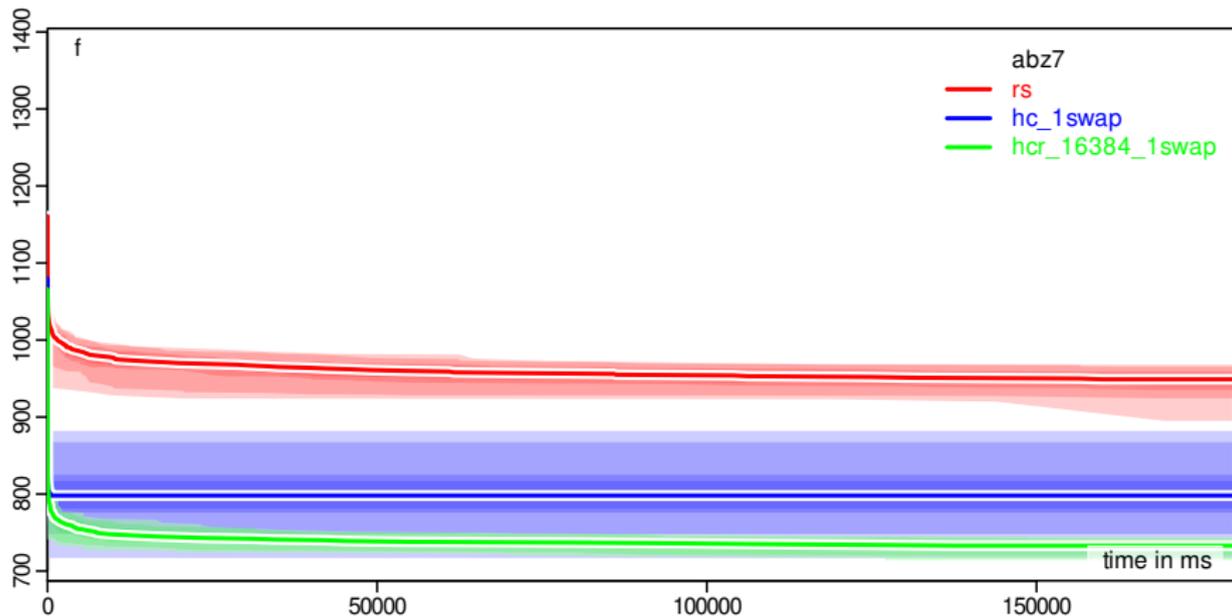
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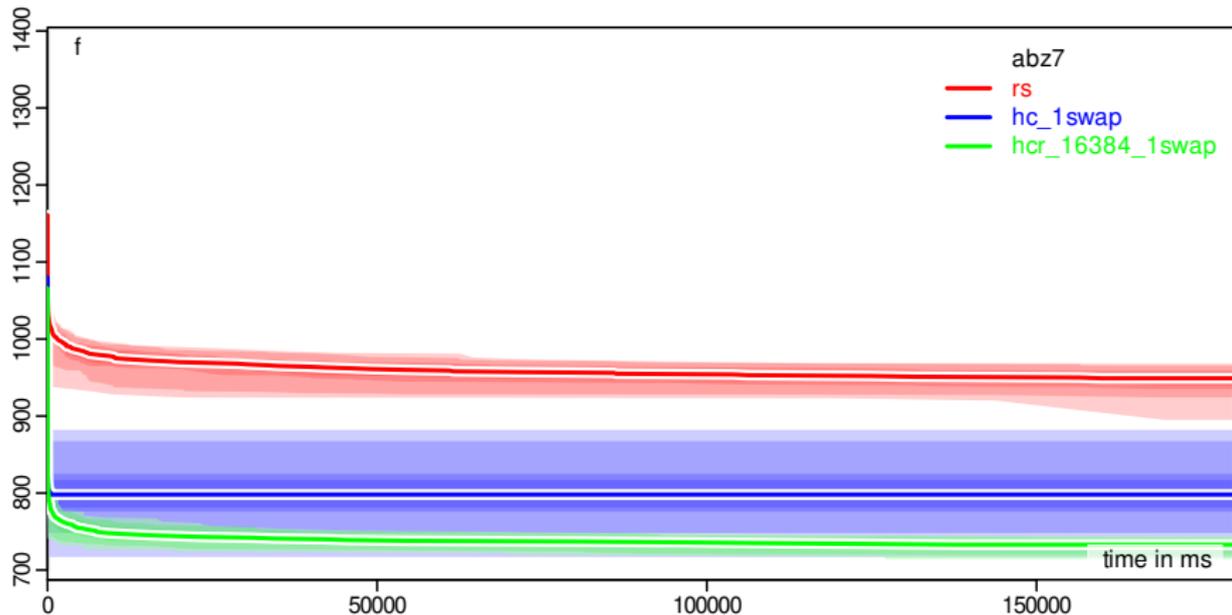
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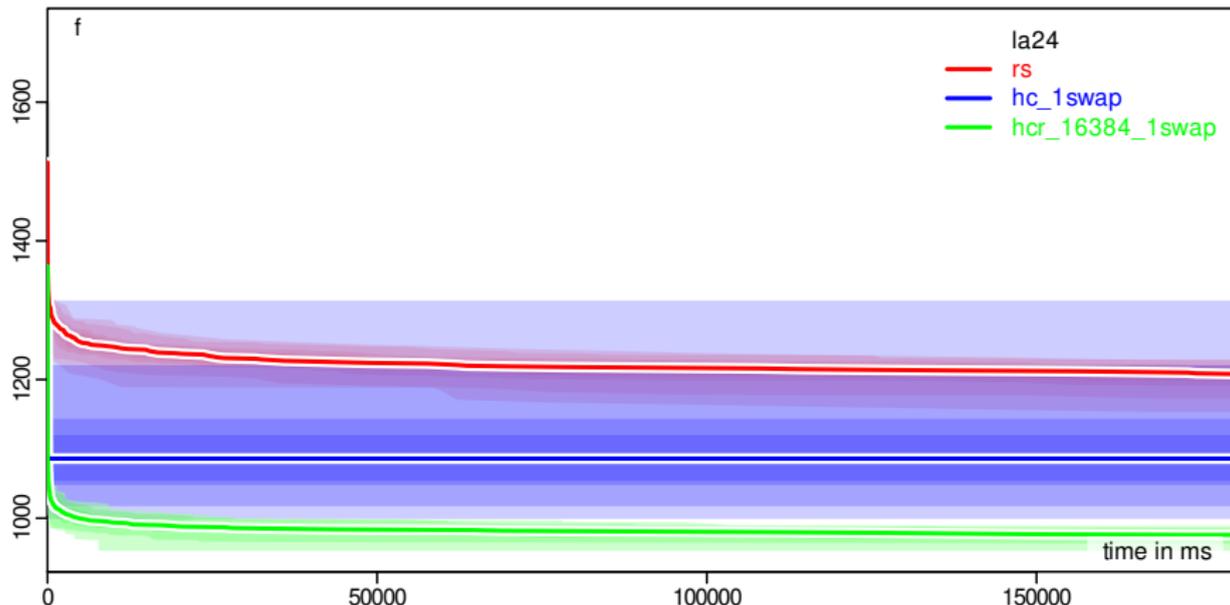
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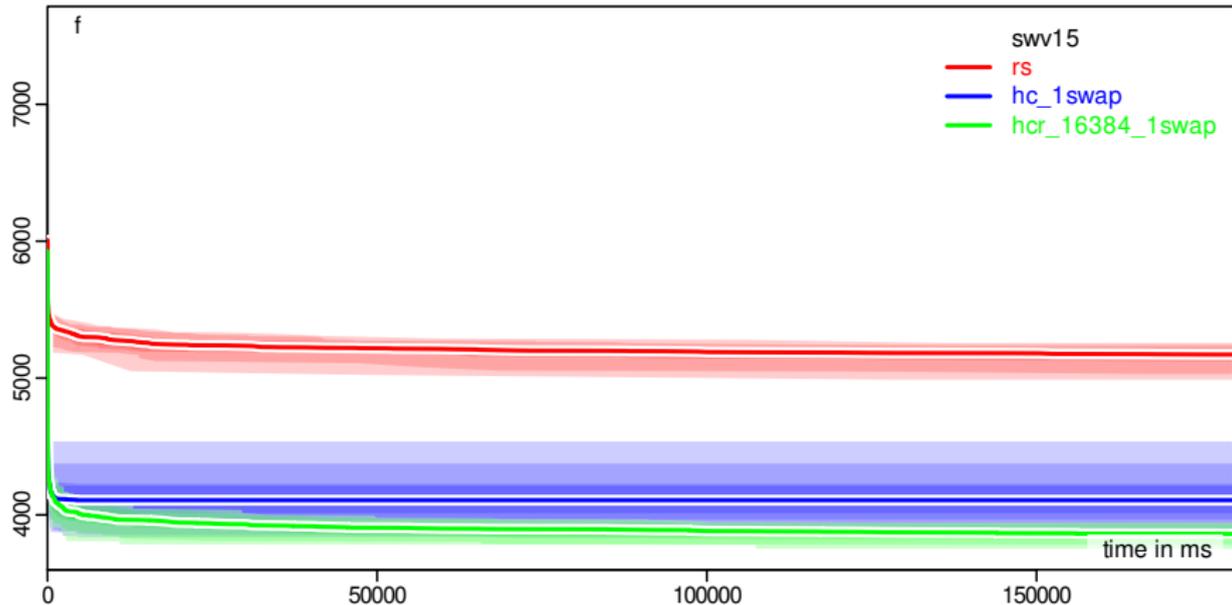
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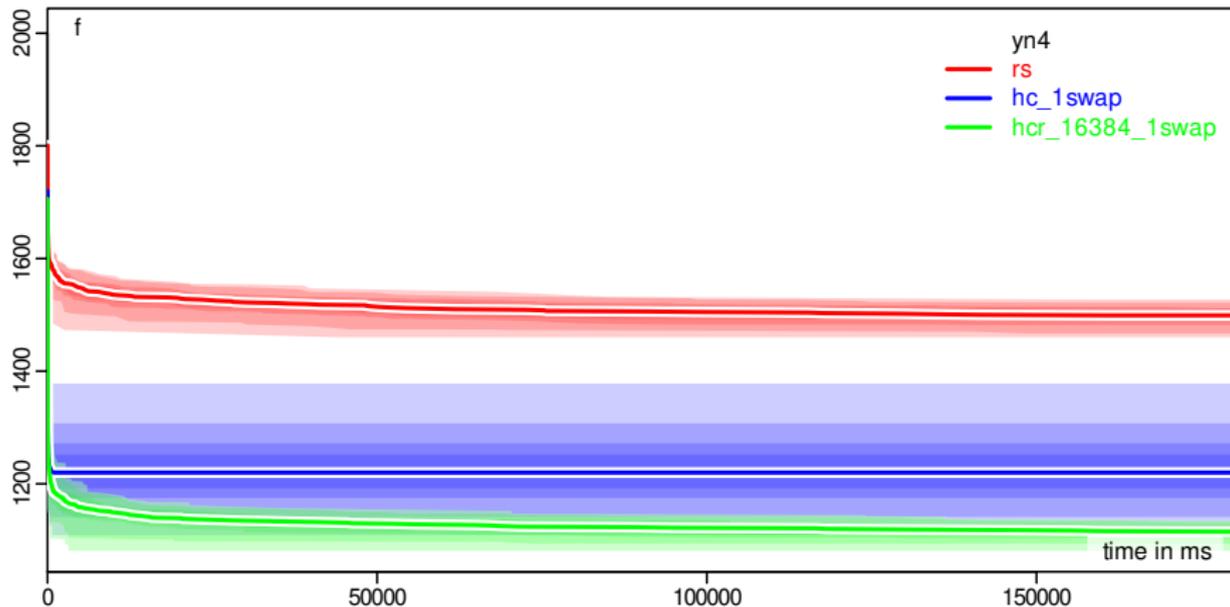
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Improved Algorithm Concept 2



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- While restarts improve the chance to find better solutions, they cannot solve the intrinsic shortcomings of an algorithm.
- Another problem is: With every restart we throw away all accumulated knowledge and information of the current run.
- Restarts are therefore also wasteful.

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- Notice: Whether or not a point x is a local optimum, is determined entirely by the unary search operator!
- If we had a different operator with a bigger neighborhood, then maybe x^\times would no longer be a local optimum and we could still improve the results after reaching it. . .

Making the neighborhood bigger

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- On the other end of the spectrum, we could simply swap all jobs in our points x randomly. Is this a good idea? Probably not: It would turn our algorithm into random sampling!
- We should respect the causality: small changes to the solution cause small changes in the objective value – big changes will lead to unpredictable results.

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Making the neighborhood bigger

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- Theoretically, we could always escape from any local optimum, but the probability may sometimes be very very small.

Implementation of the nswap Operator

```
package aitoa.examples.jssp;

public class JSSPUnaryOperatorNSwap implements IUnarySearchOperator<int[]> {
    // unnecessary stuff omitted here...
    public void apply(int[] x, int[] dest, Random random) {
        System.arraycopy(x, 0, dest, 0, x.length); // copy x to dest
        int i = random.nextInt(dest.length); // index of first job to swap
        int first = dest[i];
        //
        //
        //
        //
        //
        int j = random.nextInt(dest.length);
        int jobJ = dest[j];
        if (first != jobJ) {
            //
            //
            //
            //
            //
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            //
            //
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        inner: for (;;) { // find a location with a different job
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            int jobJ = dest[j];
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                dest[i] = jobJ; // overwrite job at index i with jobJ
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        boolean hasNext;
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So what do we get?

- I execute the program 101 times for each of the instances abz7, la24, swv15, and yn4

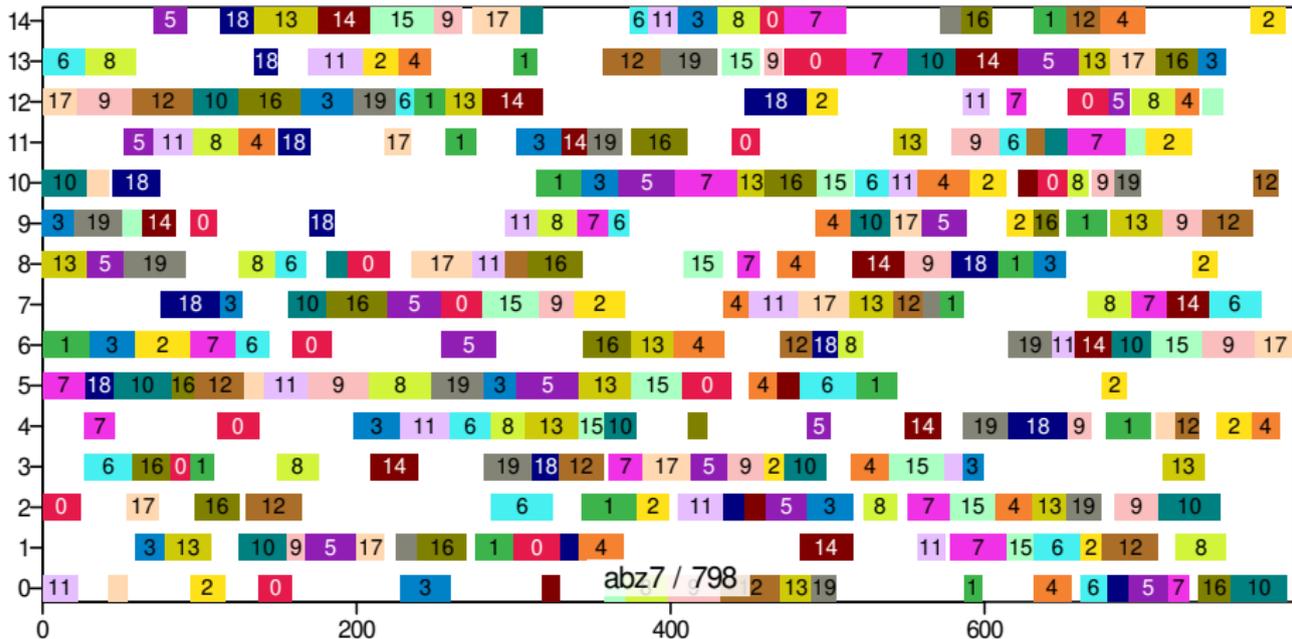
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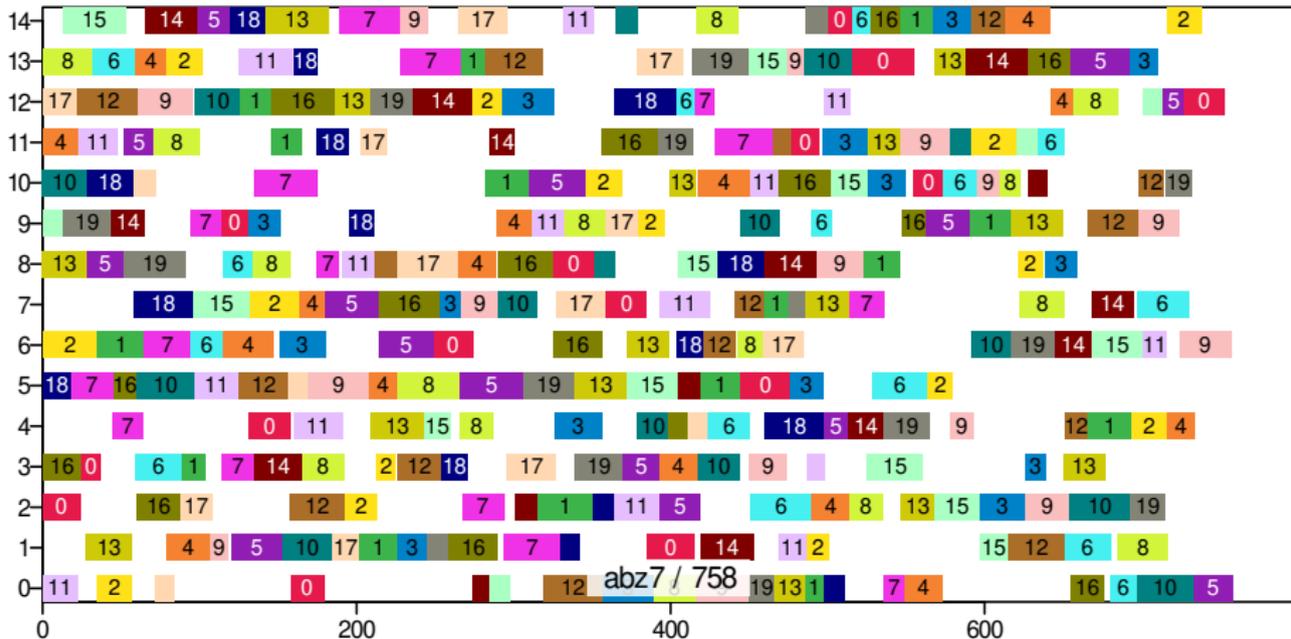
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hc_1swap: median result of 3 min of hill climber using 1swap



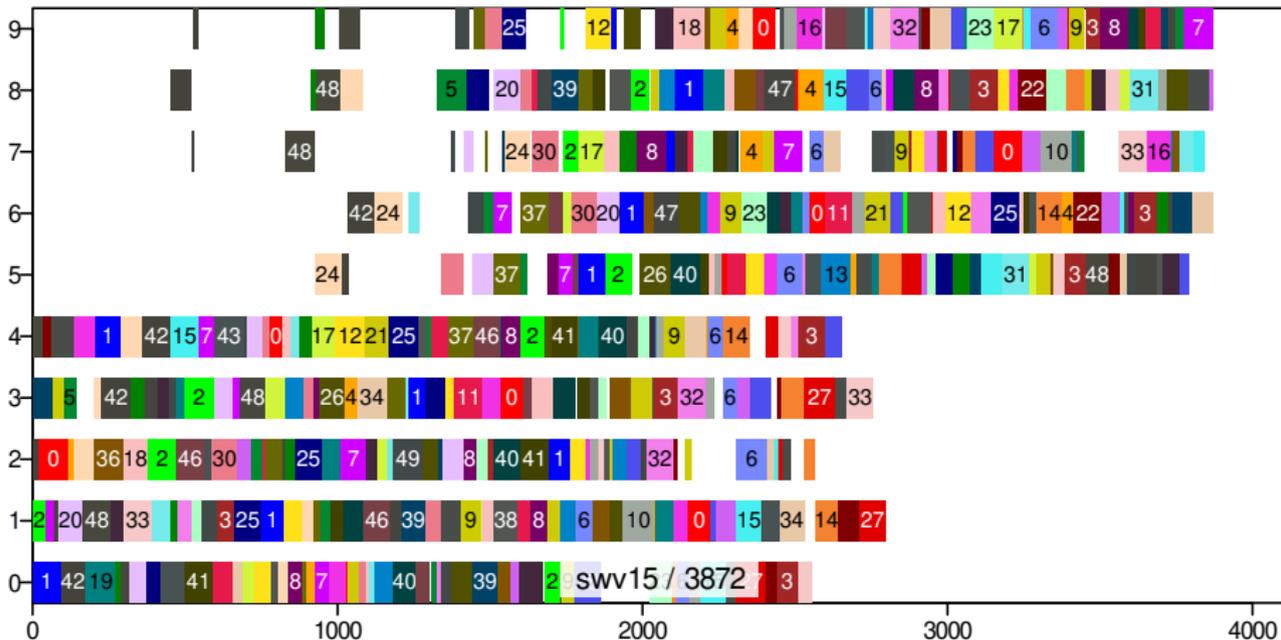
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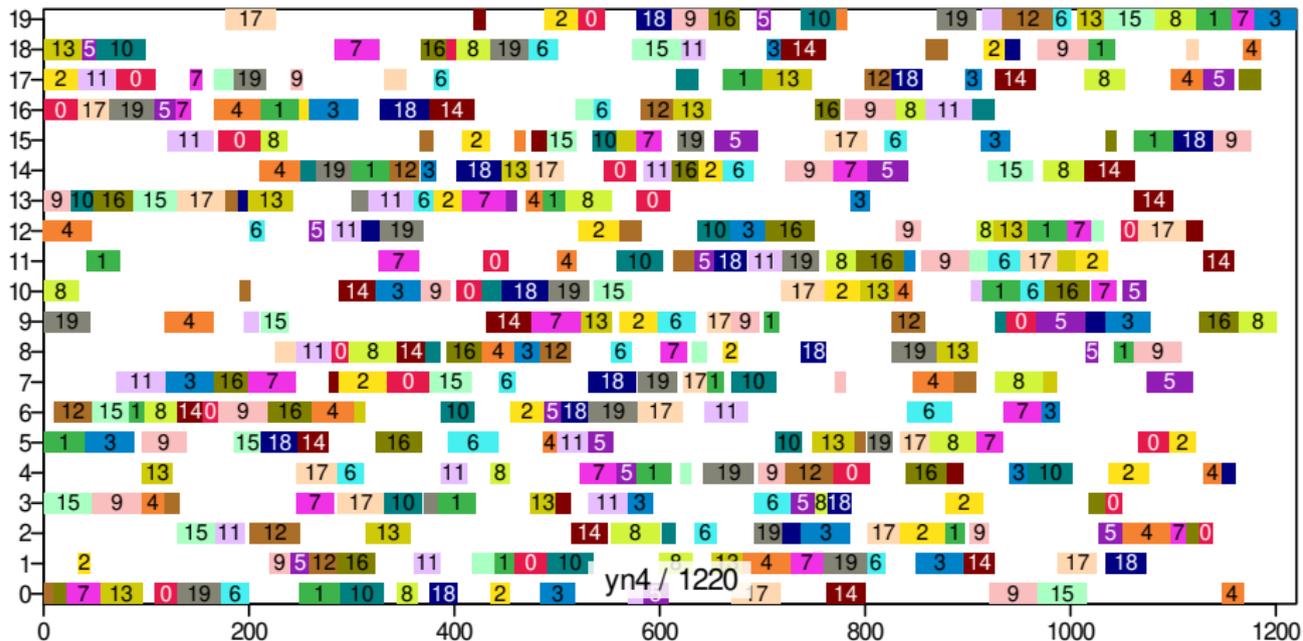
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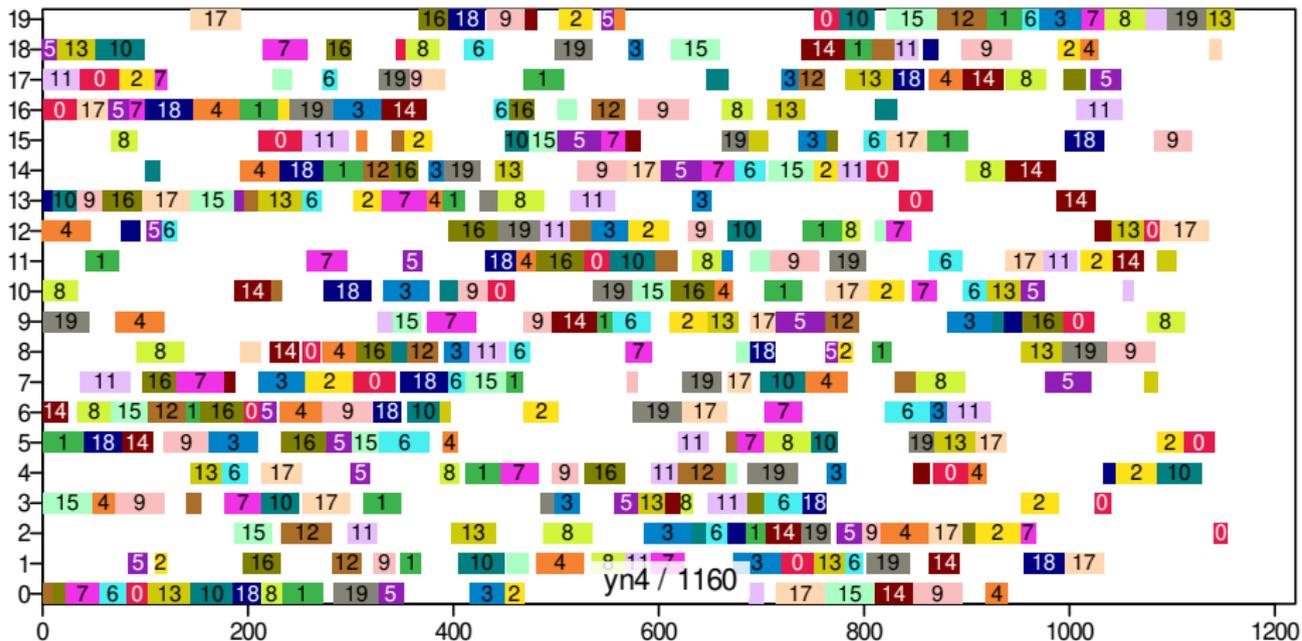
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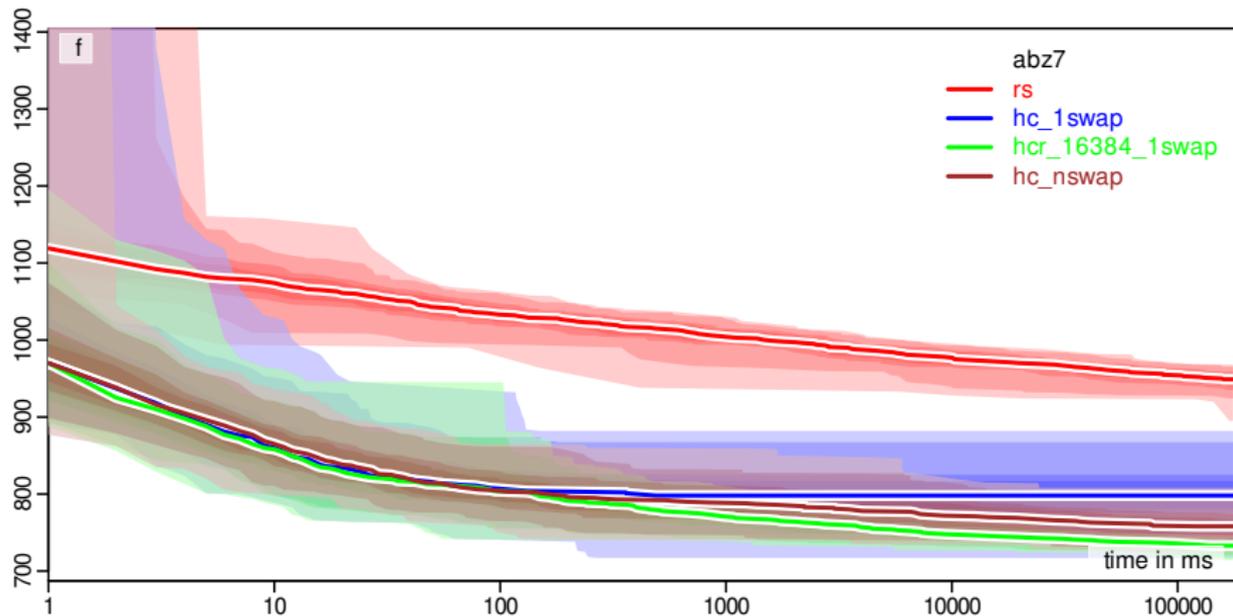
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Progress over Time

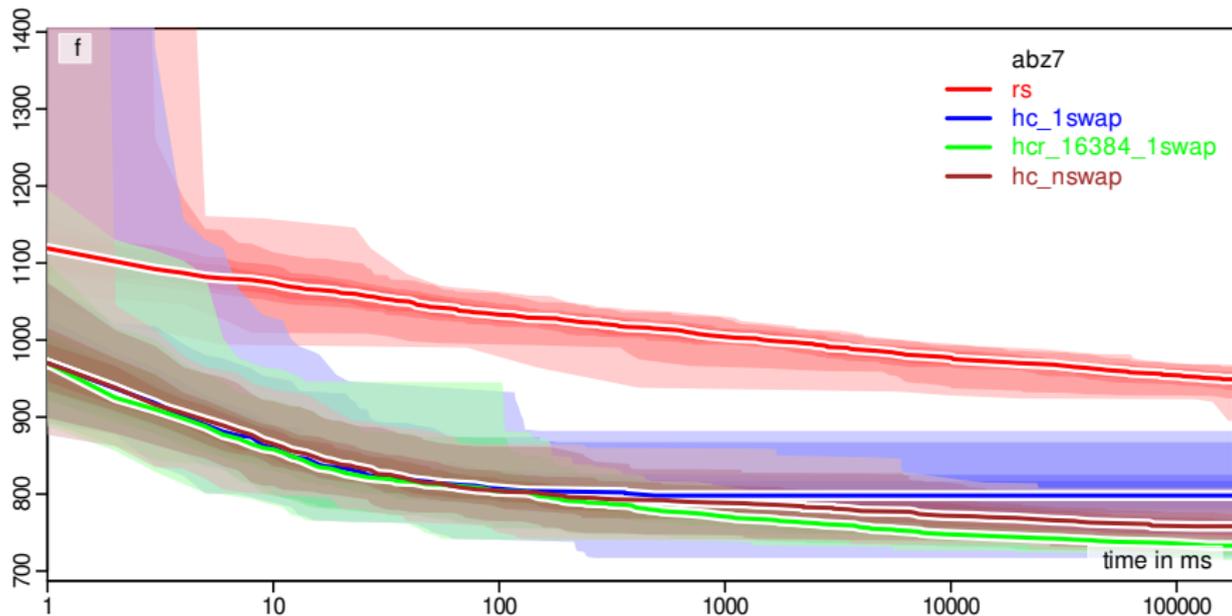
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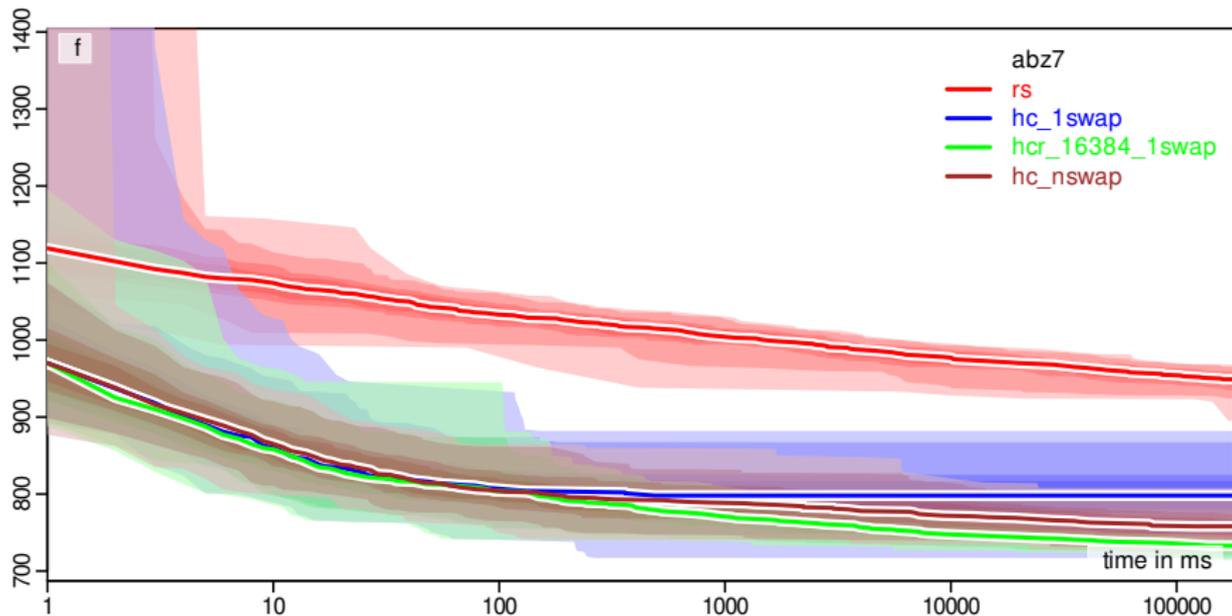
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What progress does the algorithm make over time?

- `hc_nswap` first behaves like `hc_1swap`, because most of the `nswap` moves are the same as `1swap` moves.

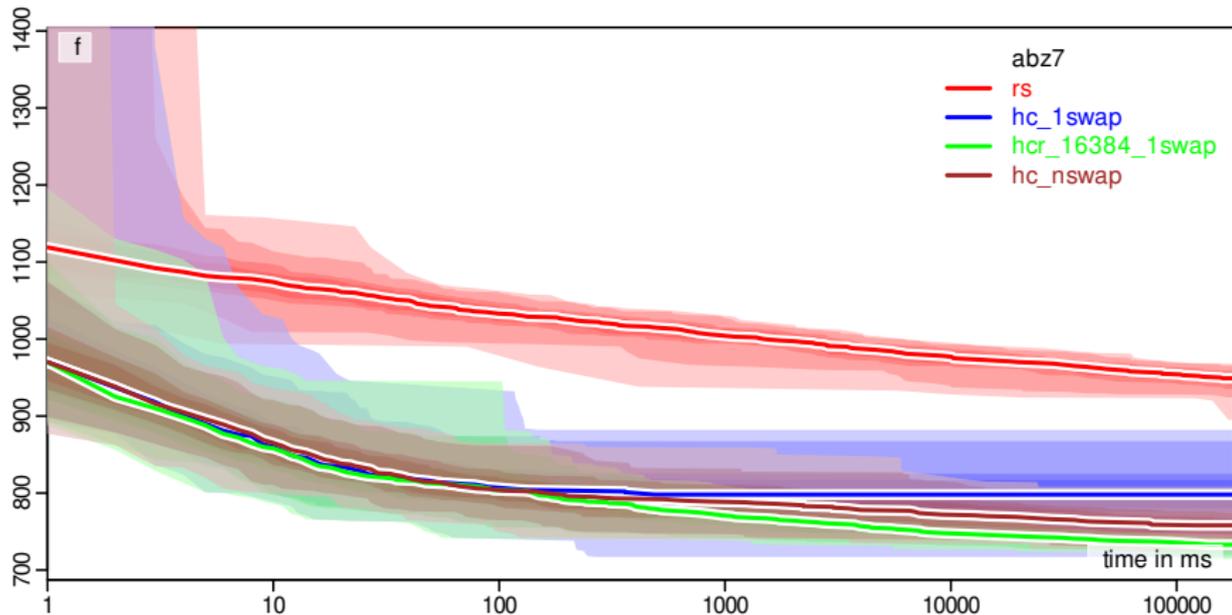
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- The rare larger moves allow it to escape from local optima that would trap `hc_1swap`.

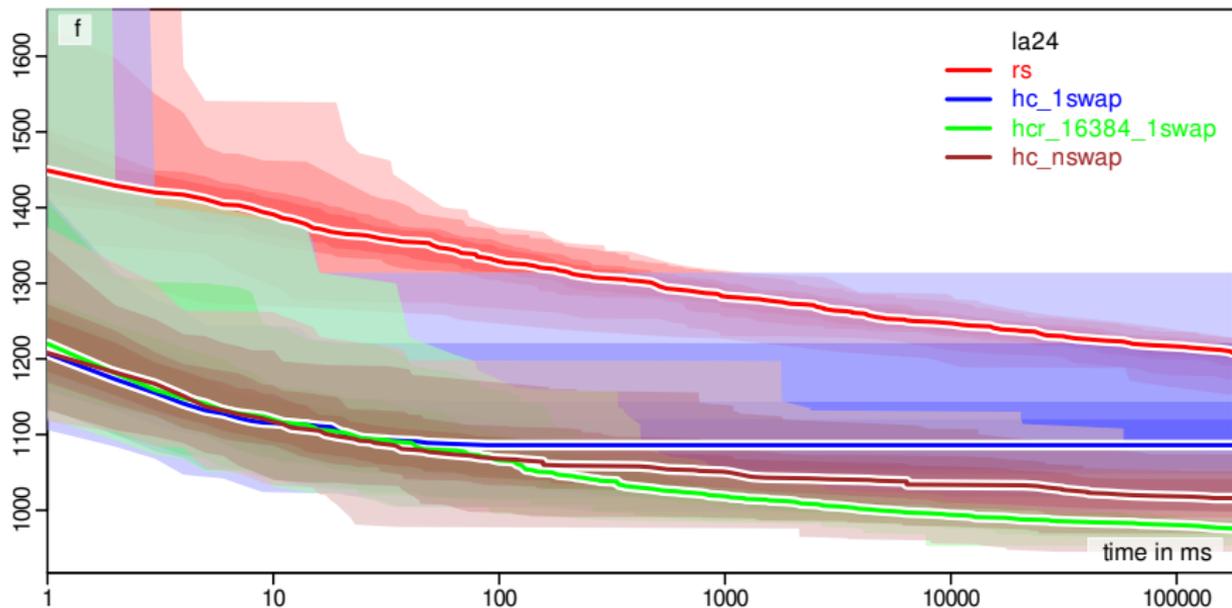
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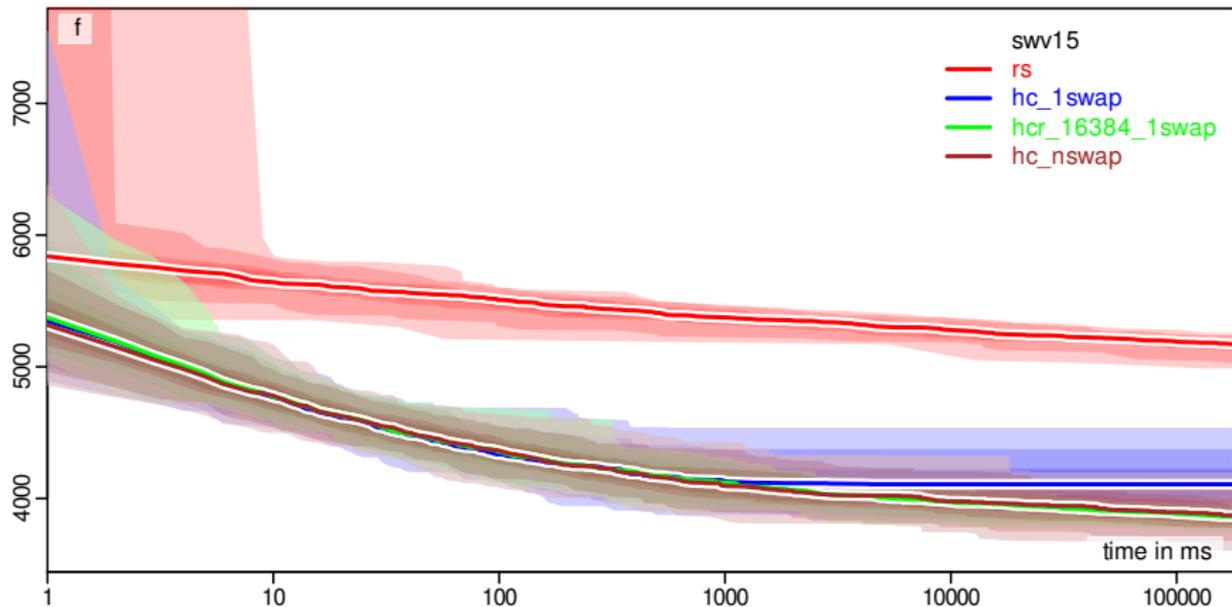
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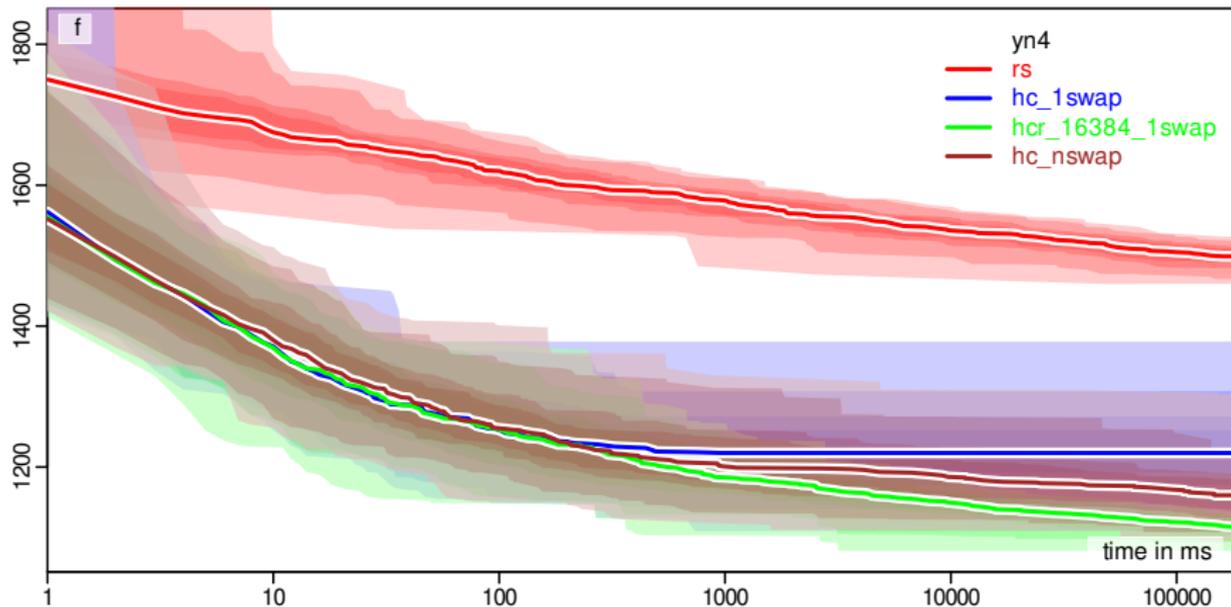
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Improved Algorithm Concept 3



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 2. we can use a unary operator with larger neighborhood that still mostly makes small steps.
- It is only natural to try to combine these two improvements.

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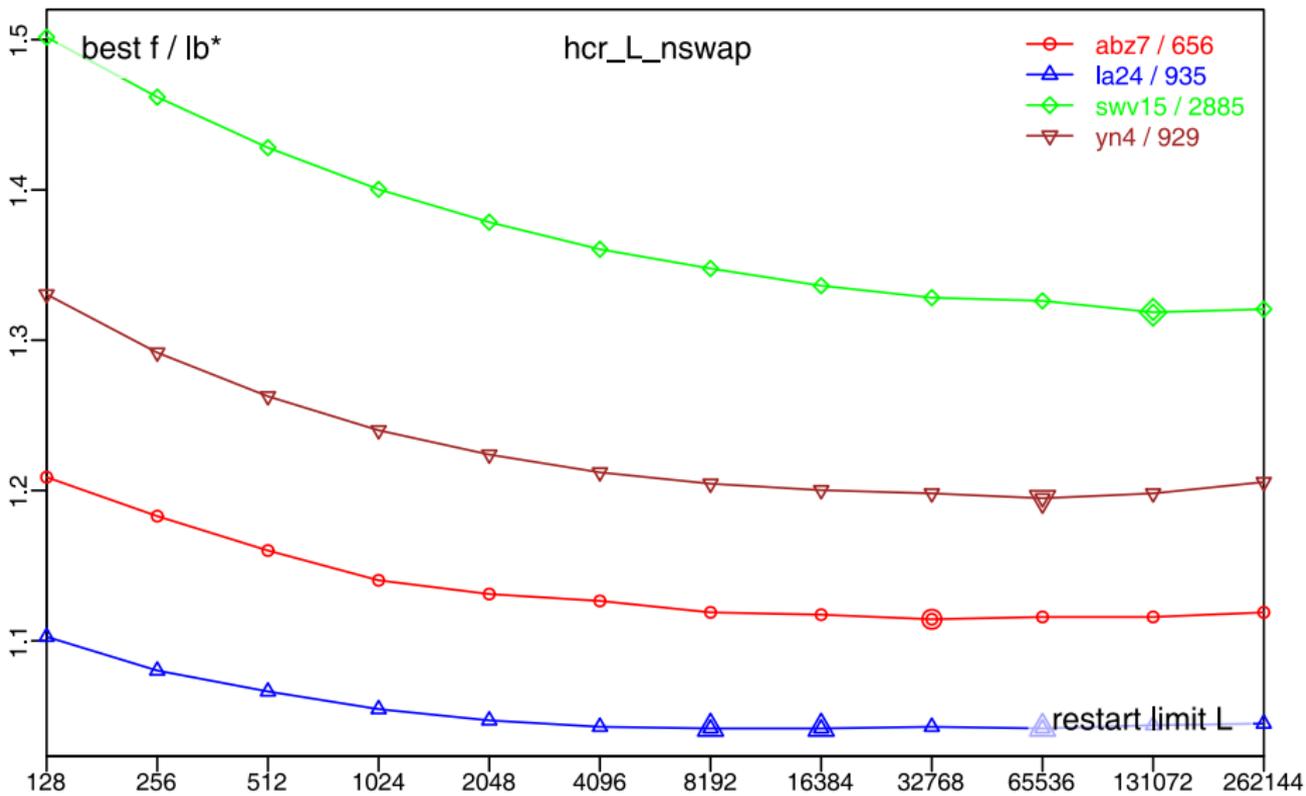
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- Let's choose $L = 65'536$, i.e., `hcr_65536_nswap`.

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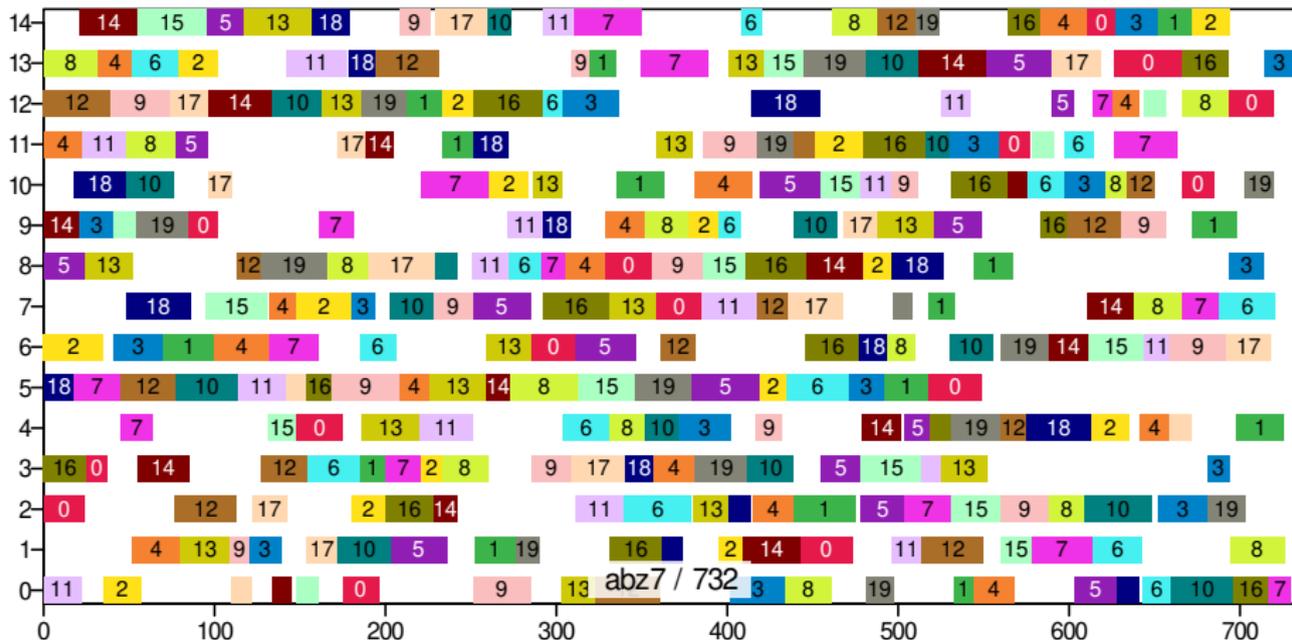
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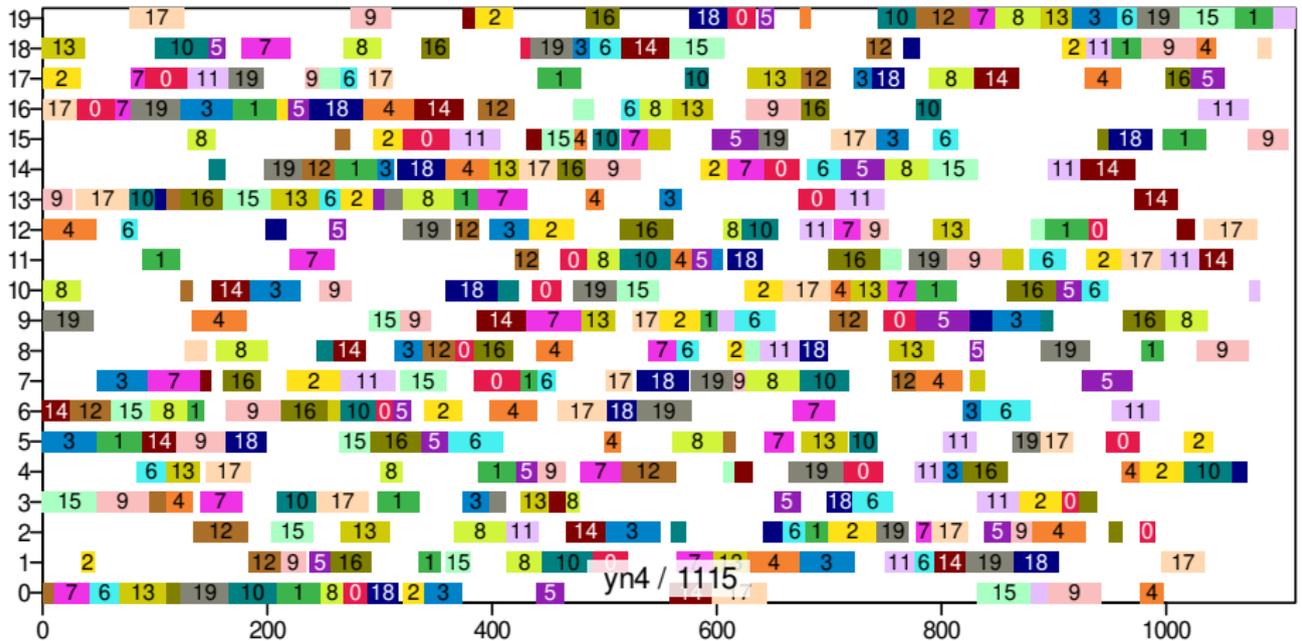
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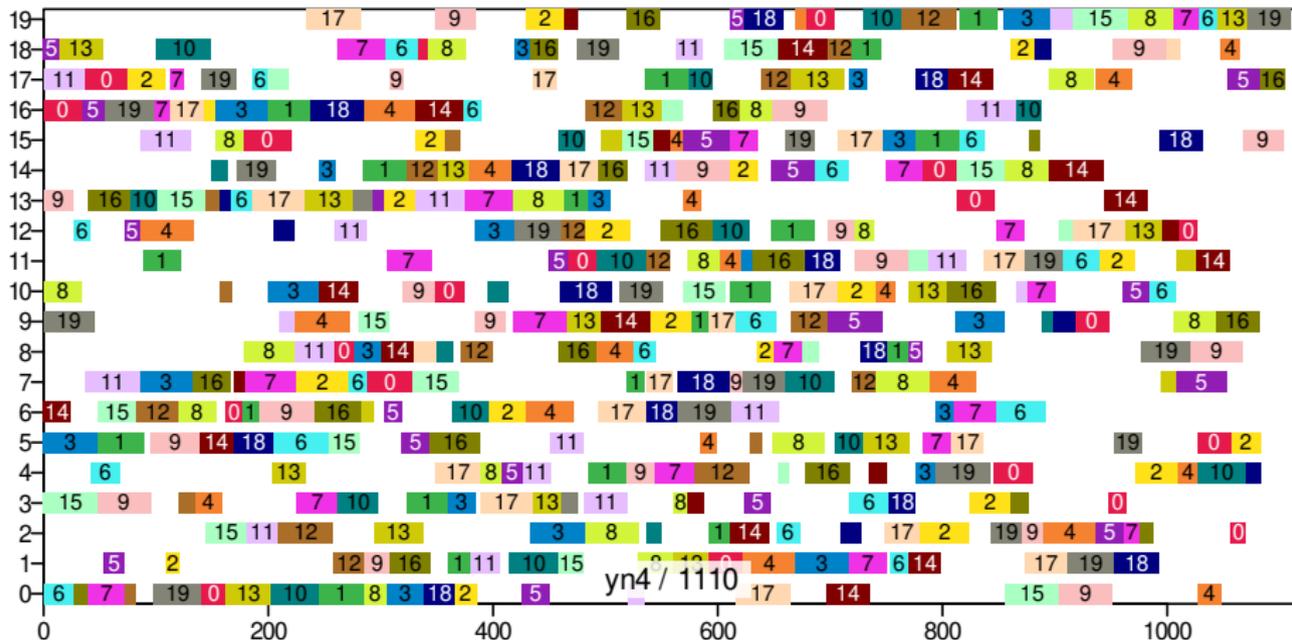
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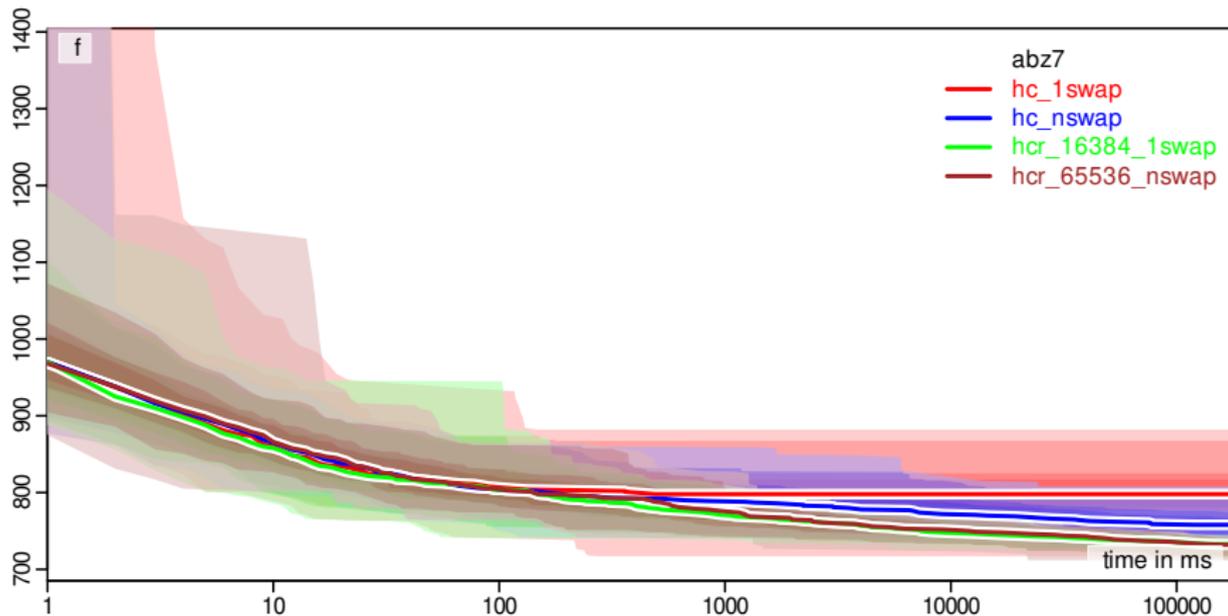


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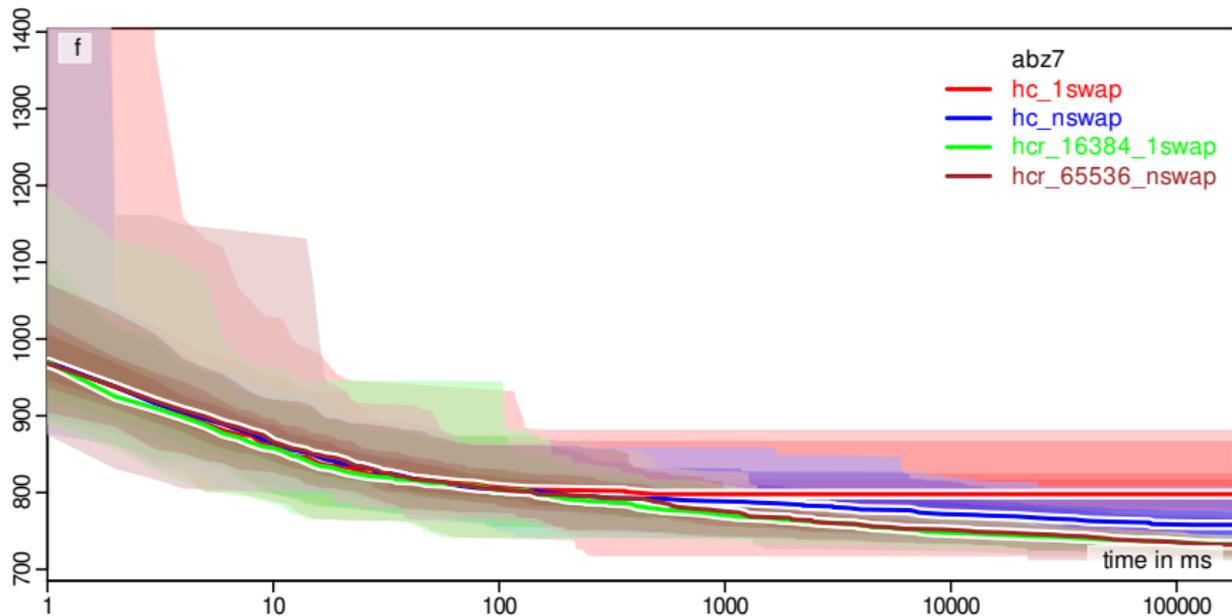
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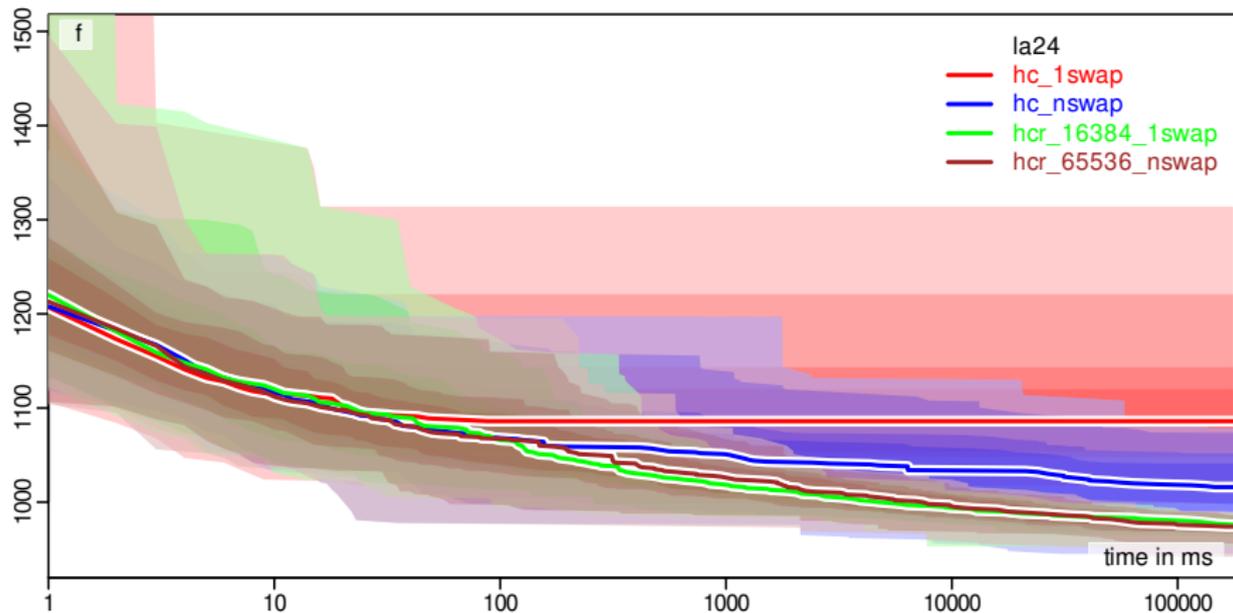
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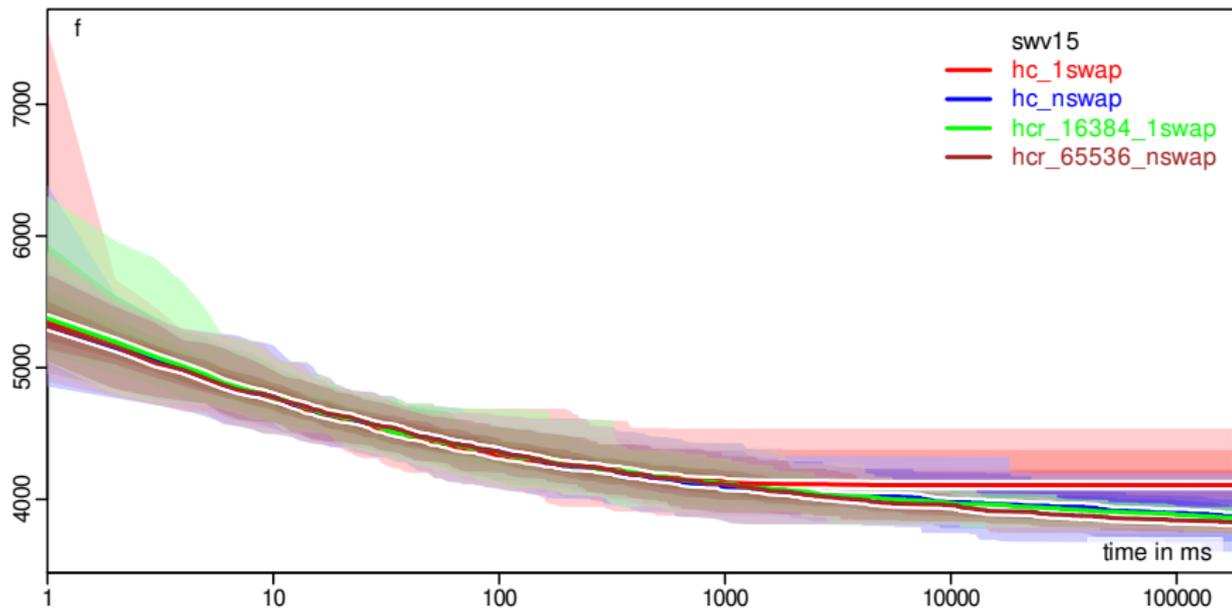
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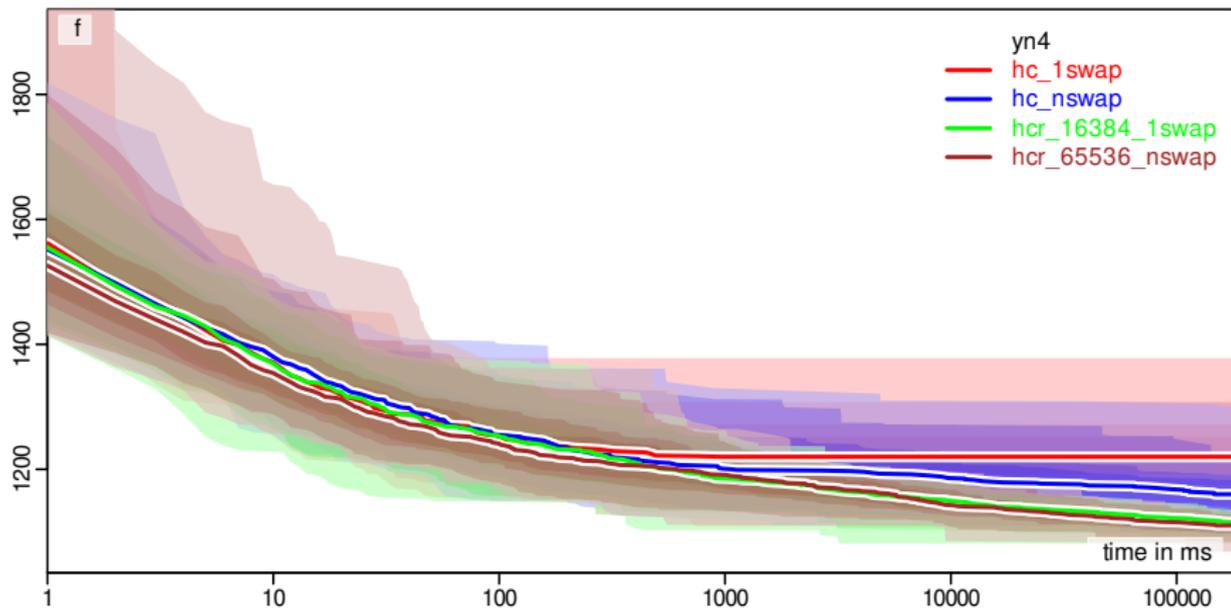
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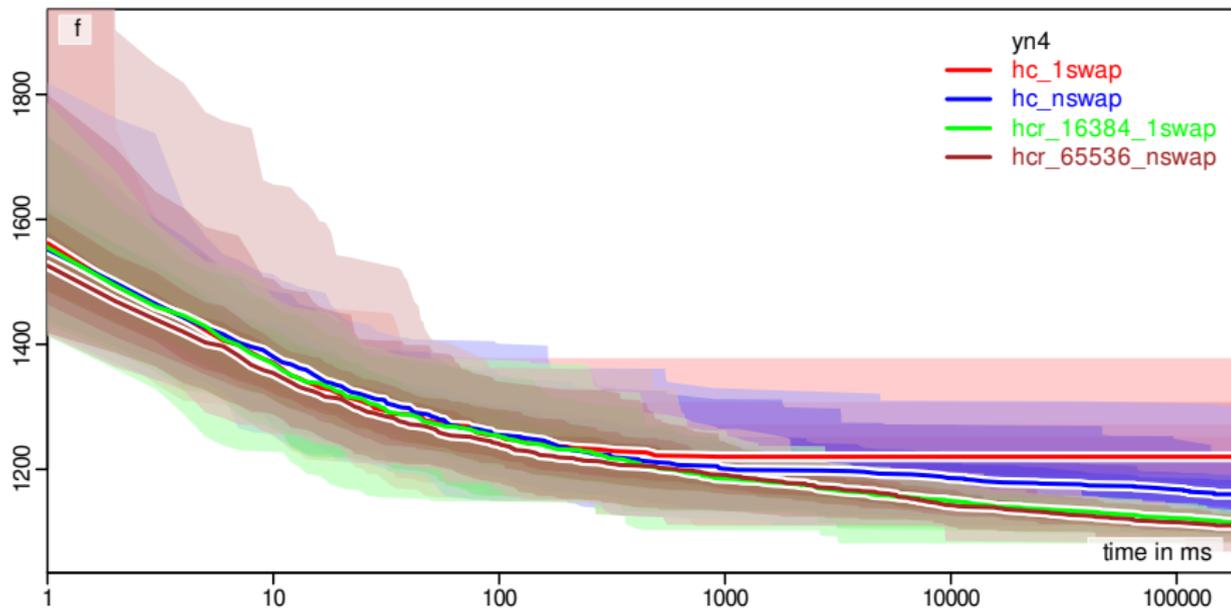
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- Hill climbing is a local search and vulnerable to get trapped in local optima.
- We can try to work around that by implementing good search operators and by restarting the algorithm.

谢谢

Thank you



References I

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